Artificial Deadlock Detection and Correction in **Bounded Scheduling of Process Networks** by **Basu Vaidyanathan EE382C - E**mbedded Software Systems **Fall 1999**

Goals

- Understand the bounded scheduling of process networks
- Develop an algorithm and implement to detect the artificial deadlock and to resolve it to continue the program
- Understand the existing basic PN framework implementation
- Modify the code to keep it modular and transparent to applications

- A networked set of Turing machines
- Models functional parallelism and simulation possible on SMP hardware
- Well-suited for signal processing systems that deal with infinite streams of data
- Termination and Boundedness are undecidable.

- Kahn process networks model:
 - has finite set of processes and FIFO queues
 - execution of a process suspended on read from an empty queue
 - a process cannot wait for data from one queue or another
 - a process may not test for presence or absence of data
 - Systems that follow Kahn's model are determinate

- Karp and Miller Computation Graph:
 - requires a threshold number of tokens on the arc before the consumer can fire
- Number of tokens produced/consumed is known only at runtime
- Dynamic scheduling is needed. It requires:
 - 1. Non-terminating programs must execute forever
 - If possible, tokens accumulation on any of the FIFO queues must be bounded

- Parks Scheduling policy has three rules:
 - 1. Process suspended when reading an empty queue
 - 2. Process suspended when writing to a full queue
 - 3. On artificial deadlock, increase the smallest full queue size until a producer can fire.
- Realizes program execution forever with bounded memory whenever possible.

Artificial Deadlock

 Occurs when atleast one process is suspended on write to a full queue

• True Deadlock

 If all the processes are suspended on read then the program has terminated

Basic Process Networks Framework

- Implementation details:
 - Developed by Greg Allen of ARL at UT
 - Implemented in C++, combined with POSIX
 Pthread library for portability
 - Threshold and PNThreshold queue layers
 - Each node as a pthread
 - FIFO queues have input and output firing thresholds
 - Threshold amount of queue data mirrored to provide address/data continuity

Basic Process Networks Framework

- Node computation time greater than thread context switch time
- POSIX condition variable used to awaken
 consumer once data is available and to
 awaken producer once space is available
- Applied in Sonar Beamforming, a real-time problem where deadlock detection is not needed
- provides a programming model for applications

My Design and Implementation

• Details:

- Variable queue size for each FIFO queue
- Maintain a list of qEntry class sorted by queue size. qEntry has Queue id, iswriteblocked stored in shared memory
- Last thread in the network before suspending itself awakens all threads suspended on write
- Only the thread with smallest queue size
 expands its queue size and continues and rest
 of the awakened threads suspended again.
- Never gets into artificial deadlock situation
- deadlock detection handled in PN queue layer

Issues and Improvements

- When expanding the queue reallocation of queue buffer is not possible
- Our PN implementation must not introduce additional deadlock violating locking hierarchy
- Use of a dedicated thread to handle deadlock
- Last thread can avoid awakening all threads suspended on write
- Searching qEntry list can be improved

Any Questions?