

Effect of Channel Estimation Error on Bit Rate Performance of Time Domain Equalizers

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Abstract—Channel equalization plays a key role in achieving high bit rates in wireline multicarrier systems. Some VDSL systems and all standardized multicarrier ADSL systems employ time domain equalization (channel shortening) and frequency domain equalization in channel equalization. In this paper, we analyze the impact of imperfect channel estimates on the bit rate performance of four time domain equalization methods. We derive a closed-form expression for the bit rate loss due to channel estimation error. We simulate the sensitivity in bit rate performance using first-generation downstream ADSL transmission. In simulation, the minimum intersymbol interference and minimum mean square error methods are relatively insensitive to channel estimation errors vs. minimum delay spread and maximum shortening signal-to-noise ratio methods.

I. INTRODUCTION

In the late 1990s, standardization of ADSL and cable modems enabled widespread deployment of high-speed data communications to home and businesses in many countries. ADSL modems use a form of multicarrier modulation known as discrete multitone (DMT) modulation. DMT modulation provides an efficient means to combat channel dispersion. The efficiency arises from the use of the fast Fourier transform (FFT) to perform multicarrier modulation and simplify frequency domain equalization. ADSL+, ADSL2, and VDSL DMT standards have been approved over the last two years.

Multicarrier modulation partitions a broadband channel into a large number of approximately independent narrowband subchannels. With the assumption that each subchannel has a flat frequency response, the total number of bits transmitted over the entire bandwidth would be the sum of the bits transmitted in each subchannel. In DMT modulation, the number of bits assigned to each subchannel is determined by the receiver based on the signal-to-noise (SNR) ratio in that subchannel only (i.e., independent of that of other subchannels). This bit loading has the potential to maximize the achieved bit rate.

Multicarrier modulation uses a cyclic prefix (CP) as a time guard band to nullify intersymbol interference (ISI) and intercarrier interference (ICI) if the channel memory is not longer than CP length. For many wireline systems, including ADSL, channel memory is generally longer than the CP length, which causes ISI and ICI. The ISI and ICI dra-

matically lower the subchannel SNR, which in turn leads to significant degradation in achievable bit rate. Reliable channel equalization is necessary to combat the severe ISI and ICI.

A conventional DMT equalizer consists of a cascade of a time-domain equalizer (TEQ), a multicarrier demodulator (FFT), and a frequency-domain equalizer (FEQ). The TEQ is an finite impulse response (FIR) filter. The FIR filter, which is in cascade with the discretized channel, shortens the cascaded channel memory to be CP length or shorter. The FEQ compensates for the phase and amplitude distortion of the shortened channel in the FFT domain by a single division per subchannel.

Many TEQ design methods have been proposed to optimize different criteria based on a training sequence. In this paper, we focus on four TEQ design methods: minimum mean squared error (MMSE) [1], [2], maximum shortening SNR (MSSNR) [3], minimum intersymbol interference (Min-ISI) [4], [5], and minimum delay spread (MDS) [6]. These four TEQ design methods have competitive bit rate performance vs. implementation complexity.

Optimum TEQ design methods often assume perfect channel knowledge and either fully rely on this knowledge to train the TEQ or use the channel knowledge to develop a fast computation method. In practical implementations, it is necessary to estimate the channel from part or all of the training sequence before one can proceed to design the TEQ. However, it is not well known how the channel estimation error affects the bit rate performance. Also, there is no explicit expression of bit rate loss is given in literature for a DMT system with channel estimation error.

In this paper, we derive a closed form model to express the bit rate loss in a DMT system due to channel estimation error. This model is generated based on previous research results of unified approach of multicarrier equalization [7] and eigenvalue perturbation theory [8]. Through computer simulations, we compare the sensitivity of bit rate performance to the channel estimation errors of the four well-known channel knowledge based equalization methods.

II. BACKGROUND

A. Bit loading in DMT

A DMT transmitter exploits an N -point Inverse FFT (IFFT) to create $N/2$ orthogonal subchannels. For large N and adequately long CP, the channel gain and noise power

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in each subchannel are flat. SNR for the i th subchannel is estimated as

$$\text{SNR}_i = \frac{E[|X_i|^2]}{E[|X_i - \hat{X}_i|^2]} \quad (1)$$

where E stands for expectation, X_i and \hat{X}_i s are the transmitted and received symbol on the i th subchannel. Number of bits assigned to i th subchannel is then determined by

$$b_i = \log_2 \left(1 + \frac{\text{SNR}_i}{\Gamma} \right) \quad (2)$$

where Γ is SNR gap for achieving Shannon channel capacity and is constant over all subchannels given the same target bit error rate for all subchannels. Bit rate of the system is calculated as

$$R = f_s * \sum_{i \in \mathcal{S}} b_i \quad (3)$$

where f_s is symbol rate and \mathcal{S} is the set of all used subchannels.

B. TEQ design

The insertion of the CP of length ν samples enables the the linear convolution of a length N DMT symbol and a channel impulse response up to length $\nu+1$ to be equal to a length N circular convolution between them. This simplifies equalization in DFT domain on a per subchannel basis, but reduces the throughput of the channel. To minimize this reduction of throughput, a TEQ is applied to reduce the overall duration of the system (channel plus equalizer) impulse response to a predefined length.

Most published time domain equalizer design methods can be unified as a maximization of a generalized Rayleigh quotient, including MMSE, MSSNR, Min-ISI and MDS [7]:

$$\mathbf{w}^{opt} = \arg \max_{\mathbf{w}} \frac{\mathbf{w}^T \mathbf{B} \mathbf{w}}{\mathbf{w}^T \mathbf{A} \mathbf{w}} \quad (4)$$

Here, \mathbf{w} is a $L_w \times 1$ vector containing the TEQ coefficients. \mathbf{A} and \mathbf{B} are generally different matrices when formulated under different criteria. The solution is the generalized eigenvector of the matrix pair (\mathbf{B}, \mathbf{A}) corresponding to the largest generalized eigenvalue.

III. BIT RATE LOSS MODEL

A. General Formulation

The received signal \hat{X}_i for the i th subchannel at the FFT output can be written as

$$\hat{X}_i = \phi_i \mathbf{q}_i^H \mathbf{Y} \mathbf{w} \quad (5)$$

where \mathbf{Y} is a $N \times L_w$ Toeplitz matrix which contains the received signal at channel output

$$\mathbf{Y} = \begin{bmatrix} y_t(\nu) & \dots & y_t(\nu - L_w + 1) \\ y_t(\nu + 1) & \dots & y_t(\nu - L_w + 2) \\ \vdots & \ddots & \vdots \\ y_t(N + \nu - 1) & \dots & y_t(N + \nu - L_w) \end{bmatrix}, \quad (6)$$

\mathbf{q}_i^H is the i th row of DFT matrix and ϕ_i is the one tap i th FEQ. Follow the approaches provided in [9], [10] by choosing an unbiased zero forcing FEQ, we could have

$$\phi_i = \frac{E[|X_i|^2]}{E[\mathbf{q}_i^H \mathbf{Y} X_i^*] \mathbf{w}} \quad (7)$$

Substitute (5) and (7) into (1), and after some manipulations, we have

$$\text{SNR}_i = \frac{|E[\mathbf{q}_i^H \mathbf{Y} X_i^*] \mathbf{w}|^2}{\mathbf{w}^T E[\mathbf{Y}^H \mathbf{q}_i \mathbf{q}_i^H \mathbf{Y}] \mathbf{w} - |E[\mathbf{q}_i^H \mathbf{Y} X_i^*] \mathbf{w}|^2} \quad (8)$$

With this SNR model, the bit rate is actually a nonlinear function of the TEQ coefficients:

$$R = f_s * \sum_{i \in \mathcal{S}} \log_2 \left(\frac{\mathbf{w}^T \mathbf{V}_i \mathbf{w}}{\mathbf{w}^T \mathbf{U}_i \mathbf{w}} \right) \quad (9)$$

where

$$\begin{aligned} \mathbf{V}_i &= \Gamma E[|X_i|^2] E[\mathbf{Y}^H \mathbf{q}_i \mathbf{q}_i^H \mathbf{Y}] \\ &\quad + (1 - \Gamma) E[\mathbf{Y}^H \mathbf{q}_i X_i] E[X_i^* \mathbf{q}_i^H \mathbf{Y}] \\ \mathbf{U}_i &= \Gamma (E[|X_i|^2] E[\mathbf{Y}^H \mathbf{q}_i \mathbf{q}_i^H \mathbf{Y}] \\ &\quad - E[\mathbf{Y}^H \mathbf{q}_i X_i] E[X_i^* \mathbf{q}_i^H \mathbf{Y}]) \end{aligned} \quad (10)$$

For the optimum design with perfect channel knowledge, we calculate bit rate by substituting (4) into (9). However, the TEQ training usually ends up at a non-optimum $\tilde{\mathbf{w}}$ due to the presence of channel estimation error. We assume in (4), \mathbf{A} and \mathbf{B} are replaced by $\mathbf{A} + \Delta \mathbf{A}$ and $\mathbf{B} + \Delta \mathbf{B}$. Due to the different formulations of \mathbf{A} and \mathbf{B} in various channel estimation based methods, $\Delta \mathbf{A}$ and $\Delta \mathbf{B}$ are not the same error matrix in general. The generalized eigen-problem of TEQ design can be reduced to finding an eigenvector of $\mathbf{C} = \mathbf{A}^{-1} \mathbf{B}$. Though in practice we consider other approaches to solve it due to numerical implementation concerns, this approach could serve here as an analytic study of channel estimation error effects.

Suppose a $n \times n$ matrix \mathbf{C} , which has n eigenvalues λ_i s, n corresponding eigenvectors \mathbf{w}_i s and n left eigenvectors \mathbf{p}_i s, is perturbed by $\tilde{\mathbf{C}} = \mathbf{C} + \Delta \mathbf{C}$. We have

$$\begin{aligned} \tilde{\mathbf{C}} &= (\mathbf{A} + \Delta \mathbf{A})^{-1} (\mathbf{B} + \Delta \mathbf{B}) \\ &= (\mathbf{A}^{-1} - \mathbf{A}^{-1} (\Delta \mathbf{A}^{-1} + \mathbf{A}^{-1})^{-1} \mathbf{A}^{-1}) (\mathbf{B} + \Delta \mathbf{B}) \\ &= \mathbf{A}^{-1} \mathbf{B} - \mathbf{A}^{-1} (\Delta \mathbf{A}^{-1} + \mathbf{A}^{-1})^{-1} \mathbf{A}^{-1} \mathbf{B} \\ &\quad + \mathbf{A}^{-1} \Delta \mathbf{B} - \mathbf{A}^{-1} (\Delta \mathbf{A}^{-1} + \mathbf{A}^{-1})^{-1} \mathbf{A}^{-1} \Delta \mathbf{B} \\ &= \mathbf{C} + \Delta \mathbf{C} \end{aligned} \quad (11)$$

where $\Delta \mathbf{C}$ is the sum of last three items in (11)

In a practical TEQ design, usually \mathbf{C} has only one largest eigenvalue λ_k , and optimum TEQ $\mathbf{w} = \mathbf{w}_k$ in this case. Under the perturbation of $\Delta \mathbf{C}$, λ_k is replaced by $\lambda_k + \Delta \lambda_k$ and \mathbf{w}_k is changed to $\mathbf{w}_k + \Delta \mathbf{w}_k$. Write

$$\Delta \mathbf{w}_k = \sum_{i=1}^n d_i \mathbf{w}_i \quad (12)$$

where d_i are projection coefficients. We have

$$\mathbf{w}_k + \Delta \mathbf{w}_k = (1 + d_k) \mathbf{w}_k + \sum_{i \neq k} d_i \mathbf{w}_i \quad (13)$$

Since eigenvectors are determined only up to a scalar multiple, we can always set $d_k = 0$ to make $\Delta \mathbf{w}_k = \sum_{i \neq k} d_i \mathbf{w}_i$. We follow the approach in [8] to expand

$$(\mathbf{C} + \Delta \mathbf{C})(\mathbf{w}_k + \Delta \mathbf{w}_k) = (\lambda_k + \Delta \lambda_k)(\mathbf{w}_k + \Delta \mathbf{w}_k). \quad (14)$$

Using the facts

$$\mathbf{C} \mathbf{w}_i = \lambda_i \mathbf{w}_i \quad (15)$$

$$\mathbf{p}_j^H \mathbf{w}_k = 0 \text{ if } k \neq j \text{ and } \mathbf{p}_k^H \mathbf{w}_k \neq 0, \quad (16)$$

it can be shown that

$$\begin{aligned} \Delta \mathbf{w}_k &= \alpha^{-1} \sum_{k \neq i} \mathbf{w}_i \frac{\mathbf{p}_i^H (\Delta \mathbf{C}) \mathbf{w}}{(\lambda_k - \lambda_i) \mathbf{p}_i^H \mathbf{w}_i} \\ \tilde{\mathbf{w}} &= \mathbf{w} + \Delta \mathbf{w}_k = \beta \mathbf{w} \end{aligned} \quad (17)$$

where

$$\begin{aligned} \alpha &= \mathbf{I} + \sum_{k \neq i} \mathbf{w}_i \frac{\mathbf{p}_i^H (\Delta \lambda_k)}{(\lambda_k - \lambda_i) \mathbf{p}_i^H \mathbf{w}_i} \\ &\quad - \sum_{k \neq i} \mathbf{w}_i \frac{\mathbf{p}_i^H (\Delta \mathbf{C})}{(\lambda_k - \lambda_i) \mathbf{p}_i^H \mathbf{w}_i}, \end{aligned} \quad (18)$$

$$\beta = \mathbf{I} + \alpha^{-1} \sum_{k \neq i} \mathbf{w}_i \frac{\mathbf{p}_i^H (\Delta \mathbf{C})}{(\lambda_k - \lambda_i) \mathbf{p}_i^H \mathbf{w}_i}, \quad (19)$$

and \mathbf{I} is $L_w \times L_w$ identity matrix. Hence, the bit rate

$$\begin{aligned} \tilde{R} &= f_s * \sum_{i \in \mathcal{S}} \log_2 \left(\frac{\tilde{\mathbf{w}}^T \mathbf{V}_i \tilde{\mathbf{w}}}{\tilde{\mathbf{w}}^T \mathbf{U}_i \tilde{\mathbf{w}}} \right) \\ &= f_s \log_2 \prod_{i \in \mathcal{S}} \left(\frac{\tilde{\mathbf{w}}^T \mathbf{V}_i \tilde{\mathbf{w}}}{\tilde{\mathbf{w}}^T \mathbf{U}_i \tilde{\mathbf{w}}} \right) \end{aligned} \quad (20)$$

Data rate loss due to imperfect channel estimation can be written as

$$\begin{aligned} \Delta R &= f_s \left(\log_2 \prod_{i \in \mathcal{S}} \left(\frac{\mathbf{w}^T \mathbf{V}_i \mathbf{w}}{\mathbf{w}^T \mathbf{U}_i \mathbf{w}} \right) - \log_2 \prod_{i \in \mathcal{S}} \left(\frac{\tilde{\mathbf{w}}^T \mathbf{V}_i \tilde{\mathbf{w}}}{\tilde{\mathbf{w}}^T \mathbf{U}_i \tilde{\mathbf{w}}} \right) \right) \\ &= f_s \log_2 \prod_{i \in \mathcal{S}} \left(\frac{\mathbf{w}^T \mathbf{V}_i \mathbf{w} (\beta \mathbf{w})^T \mathbf{U}_i (\beta \mathbf{w})}{\mathbf{w}^T \mathbf{U}_i \mathbf{w} (\beta \mathbf{w})^T \mathbf{V}_i (\beta \mathbf{w})} \right) \end{aligned}$$

B. Case Studies

With the unified approach proposed above, we look into each different design to find out what is $\Delta \mathbf{C}$ in each case.

B.1 Maximum Shortening SNR

The MSSNR approach [3] is based solely on shortening the channel impulse response. We define channel convolution matrix as,

$$\mathbf{H} = \begin{bmatrix} h(0) & h(-1) & \dots & h(-(L_w - 1)) \\ h(1) & h(0) & \dots & h(-(L_w - 2)) \\ \vdots & \vdots & \ddots & \vdots \\ h(N-1) & h(N-2) & \dots & h(N-L_w) \end{bmatrix} \quad (21)$$

and a sliding shortening window function as

$$g(n) = \begin{cases} 1 & \text{if } \Delta \leq n \leq \Delta + \nu \\ 0 & \text{elsewhere} \end{cases} \quad (22)$$

where Δ is the transmission delay. Further, we defined

$$\mathbf{G} = \text{diag}[g(0) \quad g(1) \quad \dots \quad g(N-1)]^T \quad (23)$$

and $\mathbf{D} = \mathbf{I} - \mathbf{G}$, we have

$$\begin{aligned} \mathbf{A} &= \mathbf{H}^T \mathbf{D}^T \mathbf{D} \mathbf{H} \\ \mathbf{B} &= \mathbf{H}^T \mathbf{G}^T \mathbf{G} \mathbf{H} \end{aligned} \quad (24)$$

Channel estimation error is defined as perturbation to matrix \mathbf{H} as $\Delta \mathbf{H}$. Assume $\|\Delta \mathbf{H}\|_2 = \epsilon$ is sufficiently small, we have

$$\begin{aligned} \tilde{\mathbf{A}} &= (\mathbf{H} + \Delta \mathbf{H})^T \mathbf{D}^T \mathbf{D} (\mathbf{H} + \Delta \mathbf{H}) = \mathbf{A} + \Delta \mathbf{A} \\ \tilde{\mathbf{B}} &= (\mathbf{H} + \Delta \mathbf{H})^T \mathbf{G}^T \mathbf{G} (\mathbf{H} + \Delta \mathbf{H}) = \mathbf{B} + \Delta \mathbf{B} \end{aligned} \quad (25)$$

where

$$\begin{aligned} \Delta \mathbf{A} &= \mathbf{H}^T \mathbf{D}^T \mathbf{D} \Delta \mathbf{H} + \Delta \mathbf{H}^T \mathbf{D}^T \mathbf{D} \mathbf{H} + \mathcal{O}(\epsilon^2) \\ \Delta \mathbf{B} &= \mathbf{H}^T \mathbf{G}^T \mathbf{G} \Delta \mathbf{H} + \Delta \mathbf{H}^T \mathbf{G}^T \mathbf{G} \mathbf{H} + \mathcal{O}(\epsilon^2) \end{aligned}$$

$\Delta \mathbf{C}$ is then easily computed from \mathbf{A} , \mathbf{B} , $\Delta \mathbf{A}$, and $\Delta \mathbf{B}$. TEQ and bit rate loss computation can proceed straight forwardly.

B.2 Min-ISI

The Min-ISI method generalizes the MSSNR method by weighting the ISI in the frequency domain [4], [5], e.g., to place the ISI in unused and low SNR subchannels. Similarly, for Min-ISI, we have

$$\begin{aligned} \Delta \mathbf{A} &= \mathbf{H}^T \mathbf{D}^T \left(\sum_{i \in \mathcal{S}} \mathbf{q}_i^H \frac{S_{x,i}}{S_{n,i}} \mathbf{q}_i \right) \mathbf{D} \Delta \mathbf{H} \\ &\quad + \Delta \mathbf{H}^T \mathbf{D}^T \left(\sum_{i \in \mathcal{S}} \mathbf{q}_i^H \frac{S_{x,i}}{S_{n,i}} \mathbf{q}_i \right) \mathbf{D} \mathbf{H} + \mathcal{O}(\epsilon^2) \\ \Delta \mathbf{B} &= \mathbf{H}^T \mathbf{G}^T \mathbf{G} \Delta \mathbf{H} + \Delta \mathbf{H}^T \mathbf{G}^T \mathbf{G} \mathbf{H} + \mathcal{O}(\epsilon^2) \end{aligned}$$

B.3 MDS

MDS method [6] is to minimize so called delay spread of the effective channel impulse response. Delay spread is defined as

$$D = \sqrt{\frac{1}{h_e} \sum_{n=0}^{L_h} (n - \bar{n})^2 |h(n)|^2} \quad (26)$$

where h_e is the energy of channel impulse response, and \bar{n} is a user-defined center tap of h .

$$\begin{aligned} \Delta \mathbf{A} &= \mathbf{H}^T \mathbf{Q} \Delta \mathbf{H} + \Delta \mathbf{H}^T \mathbf{Q} \mathbf{H} + \mathcal{O}(\epsilon^2) \\ \Delta \mathbf{B} &= \mathbf{H}^T \Delta \mathbf{H} + \Delta \mathbf{H}^T \mathbf{H} + \mathcal{O}(\epsilon^2) \end{aligned}$$

where $\mathbf{Q} = \text{diag}\{[(0 - \bar{n})^2, \dots, (L_w + L_h - \bar{n})^2]\}$ is a diagonal weighting matrix.

B.4 MMSE

MMSE TEQ design [1], [2] minimizes the mean square error between the output of the physical path consisting of the channel and FIR filter and the output of a virtual path consisting of a transmission delay Δ and a target impulse response (TIR). In the case of MMSE TIR with unit norm constraint, the solution to the generalized eigenvalue problem is the optimum target impulse response \mathbf{b} .

$$\begin{aligned} \mathbf{A} &= (\Psi^T \mathbf{R}_x \Psi) - \Psi^T \mathbf{R}_x \mathbf{H} (\mathbf{H}^T \mathbf{R}_x \mathbf{H} + \mathbf{R}_n)^{-1} \mathbf{H}^T \mathbf{R}_x \Psi \\ &= [(\Psi^T \mathbf{R}_x \Psi)^{-1} + \mathbf{H} \mathbf{R}_n^{-1} \mathbf{H}^T]^{-1} \\ \mathbf{B} &= \mathbf{I}_{\nu+1} \end{aligned} \quad (27)$$

where Ψ is a $(L_h + L_w - 1) \times (\nu + 1)$ windowing matrix defined as

$$[\Psi]_{m,n} = \delta(m+n-\Delta) \begin{cases} 0 \leq m < L_w + L_h - 1 \\ 0 \leq n < \nu + 1 \end{cases} \quad (28)$$

Directly compute $\Delta \mathbf{C}$ is easier in this case,

$$\begin{aligned} \tilde{\mathbf{C}} &= (\Psi^T \mathbf{R}_x \Psi)^{-1} + (\mathbf{H} + \Delta \mathbf{H}) \mathbf{R}_n^{-1} (\mathbf{H} + \Delta \mathbf{H})^T \\ \Delta \mathbf{C} &= \mathbf{H} \mathbf{R}_n^{-1} \Delta \mathbf{H}^T + \Delta \mathbf{H} \mathbf{R}_n^{-1} \mathbf{H}^T + \mathcal{O}(\epsilon^2) \end{aligned} \quad (29)$$

Once we obtain $\tilde{\mathbf{b}} = \beta \mathbf{b}$, the TEQ $\tilde{\mathbf{w}}$ can be calculated

$$\tilde{\mathbf{w}} = \frac{((\mathbf{H} + \Delta \mathbf{H})^T \mathbf{R}_x (\mathbf{H} + \Delta \mathbf{H}) + \mathbf{R}_n)^{-1} (\mathbf{H} + \Delta \mathbf{H})^T \mathbf{R}_x \tilde{\mathbf{b}}}{\mathbf{1}^T ((\mathbf{H} + \Delta \mathbf{H})^T \mathbf{R}_x (\mathbf{H} + \Delta \mathbf{H}) + \mathbf{R}_n)^{-1} (\mathbf{H} + \Delta \mathbf{H})^T \mathbf{R}_x \tilde{\mathbf{b}}} \quad (30)$$

IV. SIMULATIONS

The simulations compare the sensitivity to bit rate performance of the different equalizer designs for a wireline communication transceiver. More specifically, we consider a downstream first generation ADSL transmission. According to the ITU ADSL standard, the IFFT and FFT lengths are 512 and the cyclic prefix length is 32. We test our designs on eight typical carrier service area (CSA) loops recommended by Bell Labs [11]. Full ADSL bandwidth is up to 1.104 MHz. A common practice in industry is to use frequency division multiplexing to allocate bi-directional transmission to different frequency bands. We adopt this approach and introduce a 5th order high pass IIR filter with passband frequency at 138 kHz to separate the downstream data from the upstream data. The signal power spectral density at the transmitter output is set equal to -40 dBm/Hz. Channel noise is modeled as an additive white Gaussian noise (AWGN) with -140 dBm/Hz power density, NEXT noise from 5 integrated services digital network (ISDN) disturbers.

Fig. 1 presents magnitude responses of the eight test loops. The average channel impulse response power is between -43 dBm and -48 dBm. In our system setup, the average received signal power P_r is around -24 dBm and noise power P_n (including crosstalk and AWGN) is about -60 dBm. We model channel estimation error as an AWGN noise with variance σ^2 . According to [12], if we adopt a commonly used frequency domain channel estimates

$$\hat{H}_i = \frac{1}{L} \sum_{k=1}^L \frac{R_{k,i}}{X_i} \quad (31)$$

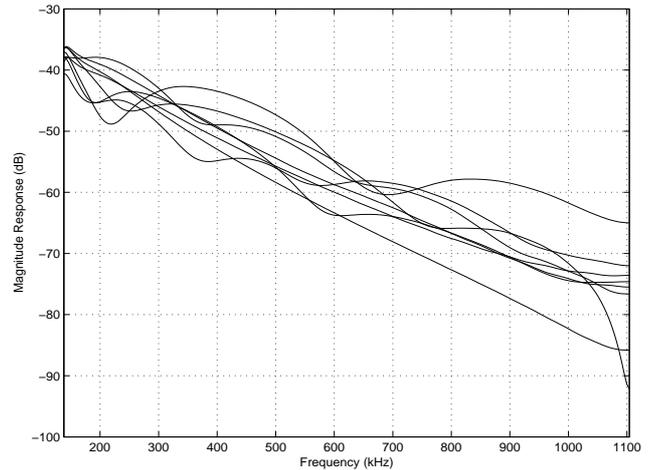


Fig. 1. Magnitude responses in downstream transmission bandwidth for eight CSA loops

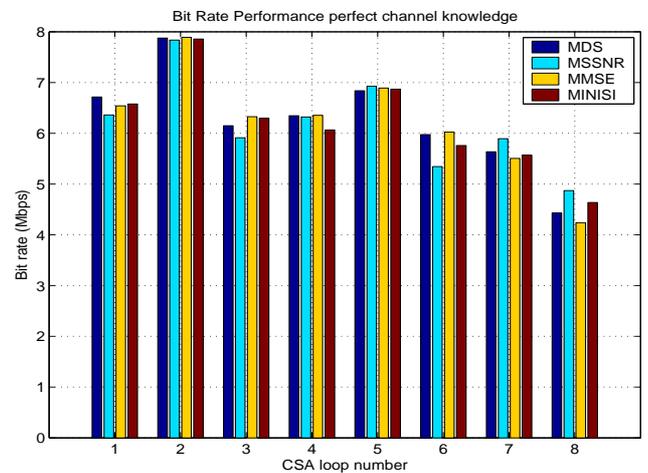


Fig. 2. Achievable bit rate for 8 CSA loops with perfect channel knowledge Coding gain is 5 dB, margin is 6 dB, input power is -40 dBm/Hz, AWGN power -140 dBm/Hz, NEXT noise is from 5 ISDN disturbers. Equalizer is trained by the MDS, MSSNR, MMSE and Min-ISI.

where $R_{k,i}$ is the i th DFT element of received channel output at k th cycle, the channel estimation error is controlled by $\sigma^2 = \frac{1}{L} P_n$. We choose a reasonable σ^2 ranging from -90 dBm to -76 dBm, where end points corresponding to averaging on $L = 1000$ cycles and $L = 40$ cycles, respectively. $L = 40$ is also suggested in [12] as a lower bound of estimation cycles. Moreover, our channel estimation error power is corresponding to an AWGN with power spectral density from -153 dBm/Hz to -133 dBm/Hz within our transmission bandwidth, which is significantly below channel gain in this range. It further suggests our choice of estimation error power is fairly conservative.

The SNR gap to Shannon capacity in our simulation is chosen as

$$\Gamma_{sim} \text{ (in dB)} = \Gamma_{gap} + \text{system margin} - \text{coding gain} \quad (32)$$

where $\Gamma_{gap} = 9.8$ dB corresponds to 10^{-7} bit error rate,

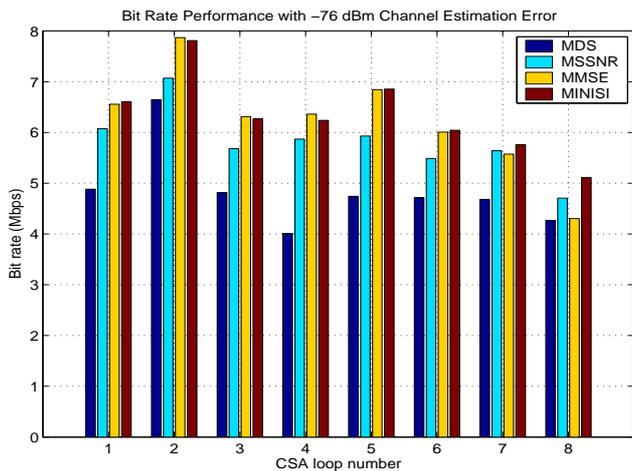


Fig. 3. Achievable bit rate for 8 CSA loops with -76 dBm channel estimation error. Coding gain is 5 dB, margin is 6 dB, input power is -40 dBm/Hz, AWGN power -140 dBm/Hz, NEXT noise is from 5 ISDN disturbers. Equalizer is trained by the MDS, MSSNR, MMSE and Min-ISI.

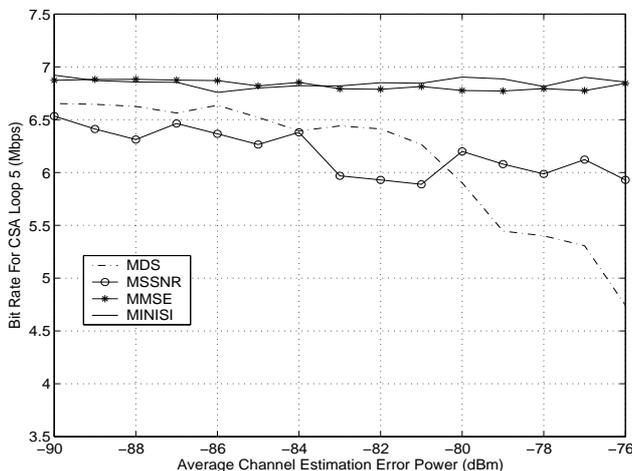


Fig. 4. Achievable bit rate for CSA loop 5 with channel estimation error from -90 dBm to -76 dBm. Coding gain is 5 dB, margin is 6 dB, input power is -40 dBm/Hz, AWGN power -140 dBm/Hz, NEXT noise is from 5 ISDN disturbers. Equalizer is trained by the MDS, MSSNR, MMSE and Min-ISI.

system margin is 6 dB, and coding gain is 5 dB.

Fig. 2 displays the achievable bit rates for eight CSA loops with perfect channel estimation. Though the four design methods use different metric to optimize TEQ settings, the bit rate performances are quite close with full knowledge of channel impulse response.

Fig. 3 displays the achievable bit rates for eight CSA loops when -76 dBm power channel estimation error is introduced. It appears that Min-ISI and MMSE outperform MSSNR by roughly 10% and MDS by roughly 20%. The performance gap is universally perceivable among all eight loops. It suggests the Min-ISI and MMSE performance hold better against channel estimation error than the performance of MSSNR and MDS.

Fig. 4 shows bit rate vs. channel estimation error power for loop 5. In Fig. 4, MSSNR and MDS are significantly affected by channel estimation error. MSSNR and MDS completely depend on the channel impulse response. In addition, MDS uses a \mathbf{Q} weighting matrix to amplify the impulse response as well as the estimation error. On the other hand, Min-ISI and MMSE have already taken noise into account, and hence, are relatively insensitive to channel estimation error. With accurate channel gain estimates, estimation error power can be 16 dB lower than additive noise power, which would likely not affect the bit rate performance as much as observed in simulations.

Since none of the four methods directly optimize the bit rate function in (2), a TEQ design with small estimation error could achieve a higher bit rate in some cases. This means ΔR of (21) might be negative in some cases.

V. CONCLUSIONS

This paper studies the behavior of bit rate performance of four popular TEQ designs when channel estimation error is present. A rate loss model as a function of optimum TEQ with perfect channel knowledge is provided. Through extensive simulations with practical noise injection, we conclude that the bit rate performance of Min-ISI and MMSE are robust against imperfect channel knowledge, while the performance of MSSNR and MDS are relatively sensitive to channel estimation error.

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