Using TestEra to Check the Intentional Naming System of Oxygen

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1. Overview

We develop TestEra (Khurshid & Marinov, 2001), a novel framework for automatic generation of test data and evaluation of correctness criteria for Java programs. As an enabling technology, TestEra uses the first-order relational language Alloy and its automatic analyzer (Jackson et al., 2001). Checking a Java program using TestEra requires: 1. Creating an Alloy model of inputs to the program; 2. Creating an Alloy model expressing the correctness criteria relating inputs and output of the program; and 3. Defining concretization and abstraction translations between Alloy model valuations and Java data structures for inputs and output. TestEra gives concrete counterexamples to violated correctness criteria. Figure 1 shows the TestEra framework.

The Intentional Naming System (INS) (Adjie-Winoto et al., 1999) is the proposed naming infrastructure for dynamic resource discovery in the MIT Oxygen project. In INS, services are referred to by *intentional* names, which describe properties that services provide. An intentional name is implemented as a tree consisting of alternating levels of *attributes* and *values*, which make up the service properties. Name resolvers in INS maintain a database that stores a mapping between service descriptions and physical network locations. Client applications invoke a resolver's *Lookup-Name* method to access services of interest. Figure 2(a) illustrates an example of invoking *Lookup-Name*.

As a case study to evaluate our framework, TestEra, we analyzed the Java implementation of INS with TestEra. Here we present the flaws TestEra identified in INS. These flaws actually existed in the INS design, and we first corrected the design. Then we fixed the code and checked its correctness using TestEra. We believe TestEra presents a novel lightweight formal method for checking Java programs.

2. Checking INS using TestEra

Our checking of INS using TestEra focuses on the *Lookup-Name* method. *Lookup-Name* returns the set of services from the input database that *conform* to the input query. To investigate the correctness of *Lookup-Name*, we test its soundness (i.e., if it returns only conforming services) and

completeness (i.e., if it returns all conforming services). The INS inventors did not state a formal definition of conformance, but only certain properties of *Lookup-Name*.

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Due to space limitation, we do not present here the Alloy model of correctness for INS, and the concretization and abstraction translations, used in TestEra's analyses. They appear elsewhere (Khurshid & Marinov, 2001). Table 1 outlines the results of the analyses explained below.

The published description of *Lookup-Name* claims: "This algorithm uses the assumption that omitted attributes correspond to wildcards; this is true for both the queries and advertisements." TestEra refutes this claim (Figure 2(b)).

TestEra also shows that addition in INS is not monotonic, i.e., addition of a new service to a database can cause existing services to erroneously become non-conforming (Figure 2(c)). This flaw points out that INS did not have a consistent notion of conformance. Both preceding flaws exist in the original design and implementation of INS.

We define a service s as conforming to a query q if s provides all the attributes and values in q in the right order (*sub-tree*). TestEra's analysis of the original implementation of *Lookup-Name* with respect to this definition of conformance reports several counterexamples. We modified the implementation and re-evaluated the correctness of *Lookup-Name* using TestEra. This time TestEra reports no flaws, increasing the confidence that our changes have corrected the problems with INS. The corrected algorithm now forms a part of the INS code base.

References

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Figure 1. TestEra framework. A TestEra model consists of Alloy and Java code. TestEra's analysis proceeds in two phases: first, all non-isomorphic instances of the Alloy input model (up to a given maximum input size) are automatically generated, and second, testing is automatically performed using appropriate concretization and abstraction translations and verification of correctness properties.



(b)

Figure 2. (a) Intentional names in INS. Query describes a camera service in building NE-43. Database stores descriptions of two services: service R0 provides a camera service in NE-43 and service R1 provides a printer service in NE-43. Invoking Lookup-Name on Query and Database should return R0. (b) TestEra's counterexample to wildcard claim. IQueryNoWC is the same as IQuery, except for the omission of the wildcarded attribute A0. Different results of the two invocations of Lookup-Name contradict the claim. (c) TestEra's counterexample to monotonicity of addition. Both services R0 and R1 are considered conforming to IQuery by the semantics of INS, but their co-existence in IDatabase makes both of them erroneously non-conforming to IQuery.

	Input size			Phase 1		Phase 2	
Property tested	Val	Att	Ser	# Tests	t[sec]	# Passed	t[sec]
Published wildcard claim	3	3	2	12	9	10 (83%)	6
Monotonicity of addition	4	2	2	160	14	150 (93%)	9
Correctness (original <i>Lookup-Name</i>)	3	3	2	16	8	10 (62%)	6
Correctness (corrected Lookup-Name)	3	3	2	16	8	16 (100%)	6

Table 1. Summary of TestEra's analyses. In all cases, TestEra takes a few seconds to complete the checking. All the properties are refuted for small input sizes. The exhaustive checking performed by TestEra uncovers subtle bugs that went undetected for over a year of use of INS.

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