Department of Electrical and Computer Engineering The University of Texas at Austin

EE 360N, Spring 2007 Yale Patt, Instructor Chang Joo Lee, Rustam Miftakhutdinov, Poorna Samanta, TAs Final Exam, May 11, 2007

Name:

Problem 1 (30 points):Problem 2 (10 points):Problem 3 (10 points):Problem 4 (15 points):Problem 5 (20 points):Problem 6 (20 points):Problem 7 (20 points):Total (125 points):

Note: Please be sure that your answers to all questions (and all supporting work that is required) are contained in the space provided.

Note: Please be sure your name is recorded on each sheet of the exam.

GOOD LUCK!

Name:_____

Problem 1 (30 points)

Part a (4 points): A RAM cell stores one bit of information. That storage is accomplished

in an SRAM cell by	
and in a DRAM cell by	

Part b (5 points): Amdahl's Law says that the maximum speedup I can get on a multiprocessor for



Part c (5 points): Recall that asynchronous I/O with central arbitration requires at least the following bus control signals:

BR_i, BG_i (for all priority levels i), BBSY, SACK, MSYN, SSYN

What signal does the PAU receive, indicating that a device controller has accepted being bus master for the next bus cycle?



What does the PAU do as a result of receiving this signal?

Problem 1 continued:

Part d (6 points): Recall that Omega networks are better than buses with respect to

and better than full crossbars with respect to



Suppose n_5 wants to talk to a_3 and n_7 wants to talk to a_4 in the following Omega network in the same cycle. Can they do it? Show how by tracing the paths on the Omega network below, or explain why not in ten words or less in the box provided.





Problem 1 continued:

Part e (5 points): A Directory based Cache Coherency scheme has five processors, each with a private cache. Cache line X has directory information as follows: 100111. Processor 2 wants to read a value in this cache line, and takes a cache miss.

What can you tell me about the Cache Coherency logic? Please be specific.

Part f (5 points): Suppose a new floating point data type is invented, which follows all the insights of the IEEE Standard. The only difference is it has 11 fraction bits and 8 exponent bits. Show below the representation for the largest positive subnormal number, using this format.

Problem 2 (10 points):

We discussed multiplication as a sequence of shifts/adds in EE 316. We extended it to Booth's algorithm (2 bits per iteration) in class. Now, on the exam, we wish to extend it to 3 bits per iteration.

A truth table to describe the combinational logic needed would consist of 4 inputs and 4 outputs. The 4 inputs are the 3 bits of the multiplier and the input borrow bit. The 4 outputs are a left shift of the multiplicand by k_1 bits, an Add or Subtract signal, a subsequent left shift of the multiplicand by k_2 bits, and the output borrow bit.

For the case where the 3 bits of the multiplier are **110** and the input borrow bit is **0**, specify the 4 outputs listed below.

1. Left shift the multiplicand

bits to the left.

2. Perform the operation of Add / Subtract (circle one) in the ALU.

3. Left shift the multiplicand

bits to the left.

4. Set the Borrow bit to 0 / 1 (circle one).

Problem 3 (10 points)

Consider the following data flow graph:



Answer:

Problem 4 (15 points):

An array A consists of 1536 one-byte elements, stored in 1536 consecutive locations of memory, starting at address x4000. (Hint : 1536 = 1024 + 512)

A programmer wishes to add 10 to every 8th element, but does so by the very curious method of adding 1 to every 8th element, and then iterating that process ten times. The C program segment looks like this:

Assume that the compiler allocates registers for the loop variables i and index. The processor has a 1 KB, 2-way set associative cache. You may assume that the cache is initially empty. You may leave the answers for the following questions as fractions. Please show your work.

Part a (5 points): For a cache line size of 16 bytes, calculate the miss ratio if the perfect LRU replacement policy is used.

Miss ratio:	
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Part b (10 points): Suppose the cache line is 4 bytes. All other parameters remain the same. Circle the replacement policy which would give the least number of misses. Also, calculate the miss ratio for your chosen replacement policy.

FIFO (First In First Out)

LIFO (Last In First Out)

LRU

Miss ratio:

Problem 5 (20 points)

An in-order vector processor with 11 cycle memory latency, 16-way interleaved memory and 8 vector registers of length 64 is used to execute the vector code resulting from compiling the following high-level language program segment:

for (i = 0; i < 40; i = i + 2)
{
 C[i]= A[i] + B[i];
 C[i+1] = A[i+1];
}</pre>

Part a (5 points): Shown below is the vector code produce by the compiler with some information missing. Your job is to fill in the missing information.

MOVI VLEN,	#							
MOVI VSTR,	#							
VLD V0, A								
VLD V1, B								
VADD V2, V0, V1								
VLD								

You have available the following vector instructions:

MOVI VLEN, #n; (1 cycle)MOVI VSTR, #n; (1 cycle)VADD Vi, Vj, Vk; Vi is dest, Vj, Vk are sources (4 staged pipelined adder)VLD Vi, A + k; Vi gets loaded with contents of memory, starting at A + kVST Vi, A + k; Contents of Vi gets stored in memory, starting at A + k

Problem 5 continued:

Part b (7 points): How many cycles does it take the vector code to execute? Assume that the memory has 2 load ports and 1 store port and the vector processor supports vector chaining. Please show your work.

Answer:

Cycles

Part c (8 points): A simple reordering of the instructions will result in improved execution time. Show the new order below which results in **minimal** execution time. How many cycles does the new order take? Please show your work.

MOVI VLEN,	#
MOVI VSTR,	#

Answer: Cycl	es
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Problem 6 (20 points)

Consider the following **two-level** virtual memory system for the LC-3b. Recall that the user page table is located in the system virtual space:

Virtual Address Space:	64KB	Number of Page Frames:	32
User Space Range:	x0000 to x7FFF	Page Size:	256 bytes
System Space Range:	x8000 to xFFFF	Page Table Entry Size:	2 bytes

The system **does not** include a Translation Lookaside Buffer and a cache. Each Page Table Entry (PTE) contains valid (V), Modified (M), Reference (R) bits and Page Frame Number(PFN). The PTE format is as follows:

15				11	10				0
V	0	0	0	Μ	R	0	0	 PFN	

Part a (2 points): How many bits are needed for the address bus to the physical memory? Show the computation.

Answer:		
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Part b (18 points): The machine stopped at a breakpoint and the following state information was observed:

PC	SBR	UBR	R0	R1	R2	R3	R4	R5	R6	R7
x3010	x0000	xA000	x1234	x5000	x4000	x1000	x4233	x1299	xFEBC	xABCD

Note: SBR is the System Page Table Base Register and UBR is the User Page Table Base Register. Each points to the first entry of the corresponding page table.

After execution resumed, the machine issued the following successive seven physical memory requests, uninterrupted by any page faults, access control violations, or anything else. Note that each entry is incomplete. Your job: complete the seven entries.

Access #	РА	Access	Data	Identity of Item Being Accessed
1		Read	x8C18	
2		Read	x8410	
3			x7E82	
4				
5			x8008	
6		Write		
7			xABCD	

Note: The last column should identify what is being read specifically. For example: "The instruction LDB R0,R1,#0," "The PTE for User Virtual Page 0," "PTE for System Virtual Page 0," etc.

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ADD⁺		00	01			DR			SR1	1	0	0	0		SR2	
ADD⁺		00	01			DR			SR 1		1		ir	nm	5	
AND⁺		01	01			DR			SR1		0	0	0		SR2	
AND⁺		01	01			DR			SR1		1		ir	nm	5	
BR		00	00		n	z	р		1	1	PC	offs	et9	1		
JMP		11	00			000		В	ase	R			000	000		
JSR		01	00		1				 	PCc	offse	et 1 1				
JSRR		01	00		0	0	0	В	ase	R			000	000		
LDB⁺		00	10			DR		В	ase	R		k	off	sete	5 1	
LDW ⁺		01	10			DR		В	ase	R			offs	et6		
LEA⁺		11	10			DR			1		PC	offs	et9			
NOT⁺		10	01			DR			SR		1		1	111	1	
RET		11	00			000			111	1		 	000	000		
RTI		10	00	1				() 000	000	000	000)	[
$LSHF^{+}$		11	01			DR			SR	1	0	0	a	mo	unt	4
$RSHFL^{+}$		11	01			DR			SR	I	0	1	a	mo	unt	4
$RSHFA^{\dagger}$		11	01	1		DR			SR	1	1	1	a	mo	unt	4
STB		00	11			SR		В	ase	R		k	off	set	5 1	
STW		01	11	1		SR		В	ase	R			offs	et6		
TRAP		11	11			00	00			I	tre	apv	vec'	18		
XOR⁺		10	01			DR			SR1		0	0	0		SR2	
XOR⁺		10	01			DR			SR		1		i	nm	5	
not used		10	10													
not used		10	11							I	I		I			

Figure 1: LC-3b Instruction Encodings

Problem 7 (20 points)

We are designing a 42-bit floating point format similar to the IEEE standard discussed in class. We have decided on the sizes of the exponent and fraction fields as shown below:



However, we have not yet decided on the bias of the exponent. Please use N for the bias in your answers to parts a, b, and c below.

Part a (5 points): What is the smallest positive number represented by this format (Min)?



Part b (5 points): What is the largest positive finite number represented by this format (*Max*)?



Problem 7 continued:

Part c (5 points): In class, we defined the range of real values rounded to *Min* if the **round to nearest** rounding mode is used. The shaded range in the figure below represents the real values rounded to *Max*. Any number to the right of that range is rounded to positive infinity.



We want to choose the bias (*N*) so that for any positive finite number *x* represented exactly by our floating point format, its reciprocal (1/x) can also be represented (after rounding to nearest) as a positive finite number.



Part d (5 points): Compute the bias that will allow the results in part c to hold. NOTE: In part c you have created two inequalities in *N*. Solve those 2 inequalities.

