

Department of Electrical and Computer Engineering
The University of Texas at Austin

EE 306, Fall 2009

Yale Patt, Instructor

Aater Suleman, Chang Joo Lee, Ameya Chaudhari, Antonius Keddis, Arvind Chandrababu, Bhargavi Narayanasetty,
Eshar Ben-dor, Faruk Guvenilir, Marc Kellermann, RJ Harden, TAs

Exam 2, November 18, 2009

Name: Solution

Problem 1 (15 points): _____

Problem 2 (15 points): _____

Problem 3 (20 points): _____

Problem 4 (10 points): _____

Problem 5 (20 points): _____

Problem 6 (20 points): _____

Total (100 points): _____

Note: Please be sure that your answers to all questions (and all supporting work that is required) are contained in the space provided.

Note: Please be sure your name is recorded on each sheet of the exam.

I will not cheat on this exam.

Signature

GOOD LUCK!

Name: _____

Problem 1. (15 points):

Part a. (5 points): Two lines of assembly code contain the following:

```
.STRINGZ "000"  
.BLKW 4
```

Does the assembler generate the identical lines of machine code for both pseudo-ops? If yes, what is the machine code. If no, explain the difference (in no more than 20 words).

No. The .STRINGZ assembles into x30, x30, x30, x0
The four words in .BLKW are not specified

Part b. (5 points): In the data path, there is a MUX (PCMUX) whose output is used to load the PC. There are three inputs to that MUX. For each input, list the LC-3 instructions that use that input at some point in their instruction cycles.

Left input	Middle input	Right input
TRAP RTI (JMP)	JMP, JSR, JSRR BR, RET	All

Part c. (5 points): In two successive clock cycles, the PC is loaded. How is that possible? Explain in no more than 20 words.

The instruction which just finished is a control instruction and the next instruction was just fetched.

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Problem 2. (15 points):

Part a. (10 points): A student wrote the following assembly language program:

```
                .ORIG  x4000
0  LEA          R2, ONE
1  AND          R5, R5, #0
TWO 2  ADD          R2, R2, #2
3  STR          R5, R2, #0
4  BRzpb       TWO
5  TRAP        x25
ONE 6  .BLKW      #50
THREE .STRINGZ  "Input a character:"
FOUR .FILL      xC000
                .END
```

PC initially contains x4000.

Does the program ever halt? Circle one: YES NO

If yes, explain what makes it halt. If no, explain why not.

When The pointer in R2 reaches x8000.

What does this program do (in no more than 15 words)?

Store 0x0000 in every other memory location from x4008-x8000

Part b. (5 points): In assembling the program of part a, the assembler creates a symbol table during the assembly process. Show the symbol table in the box below.

Symbol	Address
TWO	x4002
ONE	x4006
THREE	x4038
FOUR	x404B

x4006
x32

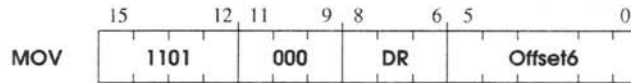
x4038

x13
x4038
x13

x404B

Name: _____

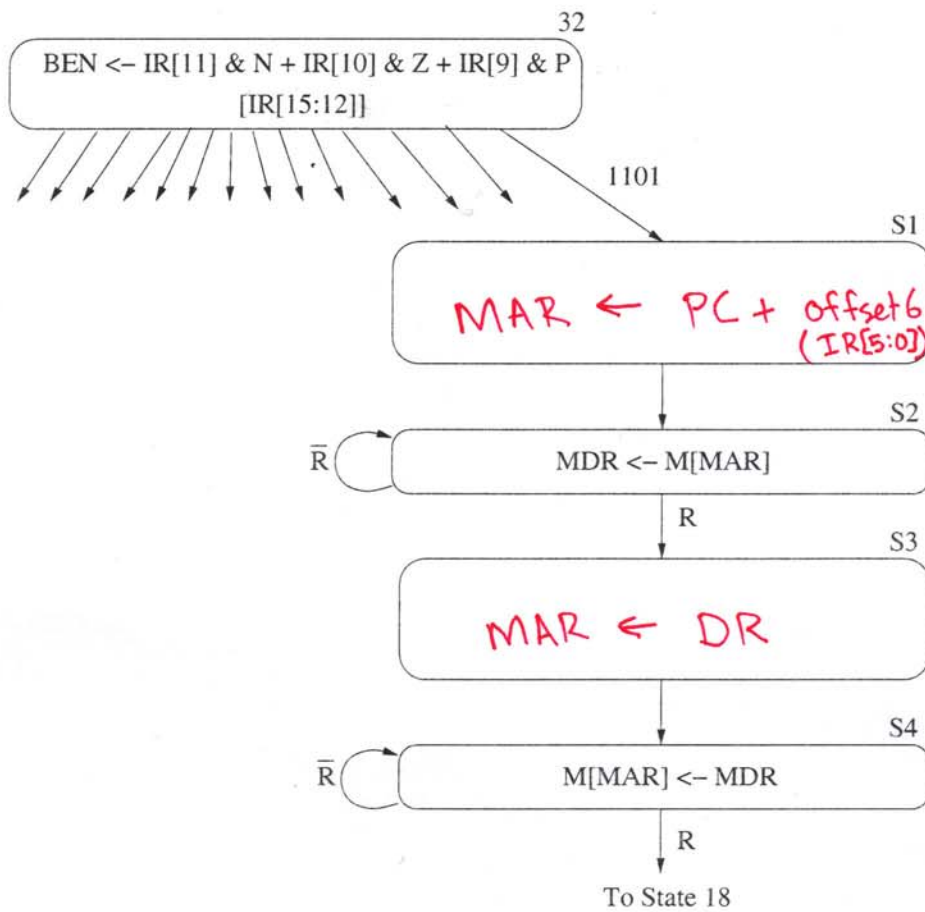
Problem 3. (20 points): As you know, in order to store into memory location A a value that is contained in memory location B, we have to first load the value into a register, and after that store it into A. Someone suggested we could do this with a single new instruction without changing the datapath in any way. Using the reserved opcode 1101, the new instruction has the following format:



The address of the source is computed by adding the incremented PC to $\text{SEXT}(\text{IR}[5:0])$. The address of the destination is in the register specified by $\text{IR}[8:6]$.

Shown below are four states of the state machine required to do this job, labeled S1, S2, S3, and S4. A partial table of control signals is shown on the next page.

Part a. (10 points): Identify what processing is necessary in states S1 and S3.



Name: _____

Part b. (10 points): Complete the table, identify the value of each control signal during each of the four states S1, S2, S3, and S4. Use the convention specified below.

Note: For a particular state, if the value of a control signal does not matter, fill it with an X.

State	ADDR1 MUX	ADDR2 MUX	MAR MUX	ALUK	LD. MAR	LD. MDR	Gate MARMUX	Gate ALU	Gate MDR	MIO.EN	R/W
S1	1	10	1	X	1	0	1	0	0	0	0
S2	X	X	X	X	0	1	0	0	0	1	0
S3	0	11	1	X	1	0	1	0	0	0	0
S4	X	X	X	X	0	0	0	0	0	1	1

- ADDR1MUX 0: SR1OUT
 1: PC
- ADDR2MUX 00: IR[10:0]
 01: IR[8:0]
 10: IR[5:0]
 11: 0
- MARMUX 0: ZEXT(IR[7:0])
 1: From adder
- ALUK 00: ADD
 01: AND
 10: NOT
 11: Pass input A
- LD.MAR 0: load not enabled
 1: load enabled
- LD.MDR 0: load not enabled
 1: load enabled
- GateMARMUX 0: do not pass signal
 1: pass signal
- GateALU 0: do not pass signal
 1: pass signal
- GateMDR 0: do not pass signal
 1: pass signal
- MIO.EN 0: memory not accessed
 1: memory accessed
- R/W 0: Read
 1: Write

Name: _____

Problem 4. (10 points): The following program is supposed to take a string of lower-case letters entered from the keyboard, convert the characters to upper-case, and then display the resulting string on the screen.

```
.ORIG x3000
LEA R0, PROMPT
TRAP x22 ; print prompt

LEA R1, INPUT
AGAIN TRAP x20 ; input a character
TRAP x21 ; echo the character
ADD R2, R0, #-10 ; subtract line feed
BRz DONE
JSR CONVERT ; call convert
STR R0, R1, #0 ; save letter
ADD R1, R1, #1 ;
BRnzp AGAIN

DONE AND R0, R0, #0
STR R0, R1, #0 ; save NULL
LEA R0, INPUT
TRAP x22 ; print output string
HALT

CONVERT ST R2, SaveR2
LD R2, MASK
AND R0, R0, R2 ; convert to uppercase
LD R2, SaveR2
RET
SaveR2 .BLKW 1

PROMPT .STRINGZ "Type a string, followed by Enter:"
INPUT .BLKW 8
MASK .FILL x00DF ; NOT of 0x20
.END
```

The program was executed twice with the following results displayed on the screen:

First time:

Type the string, followed by Enter: abcd
ABCD

Second time:

Type the string, followed by Enter: introduction
INTRODUCT@DD

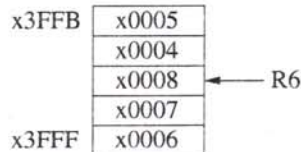
Explain why "ion" was converted into "@DD" rather than "ION". Please be specific.

The letter 'T' overwrote the MASK. Letters after 'T' are ~~AND~~ ANDed with x54, instead of xDF. For example, T(x54) AND i(x69) give @(x40). Similarly, T(x54) AND o(x6F) is D(x44)

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Problem 5. (20 points): There are times when one wants to implement a stack in memory, but can not provide enough memory to be sure there will always be plenty of space to push values on the stack. Furthermore, there are times (beyond EE 306) when it is okay to lose some of the oldest values pushed on the stack. We can save that discussion for the last class if you like.

In such situations, a reasonable technique is to specify a circular stack as shown below. In this case, the stack occupies five locations x3FFB to x3FFF. Initially, the stack is empty, with R6 = x4000. The figure shows the result of successively pushing the values 1, 2, 3, 4, 5, 6, 7, 8 on the stack.



That is, the 1 was written into x3FFF, the 2 was written into x3FFE, etc. When the time came to push the 6, the stack was full, so R6 was set to x3FFF, and the 6 was written into x3FFF, clobbering the 1 which was originally pushed.

If we now pop five elements off the stack, we get 8, 7, 6, 5, and 4, AND we have an empty stack, even though R6 contains x3FFD. Why? Because 3, 2, and 1 have been lost. That is, even though we have pushed 8 values, there can be at most only five values actually available on the stack for popping. We keep track of the number of actual values on the stack in R5.

Note that R5 and R6 are known to the calling routine, so a test for underflow can be made by the calling program using R5. Further, the calling program puts the value to be pushed in R0 before calling PUSH.

Your job: Complete the assembly language code shown below to implement the PUSH routine of the circular stack by filling in each of the four boxes with a missing instruction. We will save the POP routine for the final exam.

```
PUSH    ST R1, SAVER
        LD R1, NEGFULL
        ADD R1, R6, R1
```

BRp SKIP

```
SKIP    LD R6, BASE
        ADD R6, R6, #-1
        LD R1, MINUS5
        ADD R1, R5, R1
        BRz END
```

ADD R5, R5, #1

```
END     STR R0, R6, #0
```

LD R1, SAVER

```
RET
NEGFULL .FILL xC005 ; x-3FFB
MINUS5 .FILL xFFFB ; #-5
BASE .FILL x4000
SAVER .BLKW #1
```

Name: _____

Problem 6. (20 points): A program running on the LC-3 has reached a breakpoint. The LC-3 registers contain the following values:

PC	R0	R1	R2	R3	R4	R5	R6	R7
xF018	xFF11	xB1CD	xE000	x613C	x5151	xAC22	x1234	xF019

The computer operator immediately presses the run button, resuming execution. The table below shows a memory trace of the first seven memory accesses after execution resumes.

Hint: What state is the computer in when a breakpoint occurs? Therefore, what is the first memory access when execution is resumed?

	MAR	MDR	
1st:	x F018	x4080	; Fetched JSRR R2, PC ← R2 R7 ← xF019
2nd:	x E000	xE002	; LEA R0, #2; R0 ← xE003, PC ← PC + 1
3rd:	x E001	xBFFE	; STI R7, #-2;
4th:	x E000	x E002	; get address
5th:	x E002	x F019	; store
6th:	x E002	x F019	; TRAP x19
7th:	x 0019	x1A00	; get trap routine's address.

Your job:

Part a. (14 points): Fill in the missing values in the above table.

Part b. (6 points): What are the values in PC, R0, and R7 during the cycle in which the 8th memory access occurs.

PC: x1A00

R0: xE003

R7: xE003