Lossless Compression <u>Multimedia Systems (Module 2 Lesson 3)</u>

Summary:

Sources: The Data Compression Book, 2nd

Gailly.

- Dictionary based
 Compression
- Compression
 Adaptive Mechanism
- Limpel Ziv Welch (LZW) mechanism
- LZW Compression Article from Dr. Dobbs Journal: Implementing LZW compression using Java, by Laurence Vanhelsuwé

Ed., Mark Nelson and Jean-Loup

Dictionary-Based Compression

 The compression algorithms we studied so far use a statistical model to encode single symbols

- Compression: Encode symbols into bit strings that use fewer bits.
- Dictionary-based algorithms do not encode single symbols as variable-length bit strings; they encode variable-length strings of symbols as single tokens
 - The tokens form an index into a phrase dictionary
 - $\odot\,$ If the tokens are smaller than the phrases they replace, compression occurs.
- Dictionary-based compression is easier to understand because it uses a strategy that programmers are familiar with-> using indexes into databases to retrieve information from large amounts of storage.
 - Telephone numbers
 - Postal codes

Dictionary-Based Compression: Example

 Consider the Random House Dictionary of the English Language, Second edition, Unabridged. Using this dictionary, the string:

A good example of how dictionary based compression works can be coded as:

1/1 822/3 674/4 1343/60 928/75 550/32 173/46 421/2

- Coding:
 - Uses the dictionary as a simple lookup table
 - Each word is coded as x/y, where, x gives the page in the dictionary and y gives the number of the word on that page.
 - The dictionary has 2,200 pages with less than 256 entries per page: Therefore x requires 12 bits and y requires 8 bits, i.e., 20 bits per word (2.5 bytes per word).
 - Using ASCII coding the above string requires 48 bytes, whereas our encoding requires only 20 (<-2.5 * 8) bytes: 50% compression.

Adaptive Dictionary-based Compression

- Build the dictionary adaptively
 - Necessary when the source data is not plain text, say audio or video datá.
 - Is better tailored to the specific source.
- Original methods due to Ziv and Lempel in 1977 (LZ77) and 1978 (LZ78). Terry Welch improved the scheme in 1984 (called LZW compression). It is used in, UNIX compress, and, GIF.
- LZ77: A sliding window technique in which the dictionary consists of a set of fixed length phrases found in a window into the previously processed text
- LZ78: Instead of using fixed-length phrases from a window into the text, it builds phrases up one symbol at a time, adding a new symbol to an existing phrase when a match occurs.

LZW Algorithm

Preliminaries:

- □ A dictionary that is indexed by "codes" is used.
- $\hfill\square$ The dictionary is assumed to be initialized with 256 entries (indexed with ASCII codes 0 through 255) representing the ASCII table.
- □ The compression algorithm assumes that the output is either a file or a communication channel. The input being a file or buffer.
- Conversely, the decompression algorithm assumes that the input is a file or a communication channel and the output is a file or a buffer.



LZW Algorithm



character.

| put String: ^wed^we^wee^web^wet | | | | | | |
|---------------------------------|---|--------|-------|--------|--|--|
| w | k | Output | Index | Symbol | | |
| NIL | ^ | | | | | |
| ^ | w | ^ | 256 | ^w | | |
| w | E | w | 257 | WE | | |
| E | D | E | 258 | ED | | |
| D | ^ | D | 259 | D^ | | |
| ^ | w | | | | | |
| ^w | E | 256 | 260 | ^WE | | |
| E | ^ | E | 261 | E^ | | |
| ^ | w | | | | | |
| ^W | E | | | | | |
| ^WE | E | 260 | 262 | ^WEE | | |
| E | ^ | | | | | |
| E^ | w | 261 | 263 | E^W | | |
| w | E | | | | | |
| WE | В | 257 | 264 | WEB | | |
| R | ^ | B | 265 | B^ | | |

| set w | = NIL |
|---------|--------------------------------|
| loop | |
| | read a character k |
| | if wk exists in the dictionary |
| | w = wk |
| | else |
| | output the code for w |
| | add wk to the dictionary |
| | w = k |
| endloop | |
| | |
| | |
| | |

LZW Algorithm

LZW Decompression:

w

Е

т

EOF

^W

^WE

т

read fixed length token k (code or char) output k w = k loop read a fixed length token k

260

Т

266

^WET

entry = dictionary entry for k output entry add w + first char of entry to the dictionary

w = entry

endloop

The nice thing is that the decompressor builds its own dictionary on its side, that matches exactly the compressor's, so that only the codes need to be sent.

Example of LZW

Input String (to decode): ^WEDx256>E<260x261>257>B<260>T

| w | k | Output | Index | Symbol |
|-----|-------|--------|-------|--------|
| | ^ | ^ | | |
| ^ | W | W | 256 | ^W |
| W | E | E | 257 | WE |
| E | D | D | 258 | ED |
| D | <256> | ^W | 259 | D^ |
| ^W | E | E | 260 | ^WE |
| E | <260> | ^WE | 261 | E^ |
| ^WE | <261> | E^ | 262 | ^WEE |
| E^ | <257> | WE | 263 | E^W |
| WE | В | В | 264 | WEB |
| В | <260> | ^WE | 265 | Β^ |
| ^WE | Т | Т | 266 | ^WET |

read a fixed length token k (code or char) output k w = k loop read a fixed length token k (code or char) entry = dictionary entry for k output entry add w + first char of entry to the dictionary w = entry endloop

| LZW Algorithm - Discussion | | | | | | | | | | |
|--|--|--|--|--|--|--|--|--|--|--|
| Where is the compression? | | | | | | | | | | |
| Original String to decode : ^WED^WE^WEE^WEB^WET | | | | | | | | | | |
| Decoded String : ^WED<256>E<260><261><257>B<260>T | | | | | | | | | | |
| Plain ASCII coding of the string : 19 * 8 bits = 152 bits | | | | | | | | | | |
| LZW coding of the string: 12*9 bits = 108 bits (7 symbols and 5 | | | | | | | | | | |
| coaes, each of y dits) | | | | | | | | | | |
| • An ASCII character has a value ranaina from 0 to 255 | | | | | | | | | | |
| • All tokens have fixed length | | | | | | | | | | |
| There has to be a distinction in representation between an ASCII character and a Code (assigned to strings of length 2 or | | | | | | | | | | |
| more) | | | | | | | | | | |
| Codes can only have values 256 and above | Codes can only have values 256 and above | | | | | | | | | |
| · · · · · · · · · · · · · · · · · · · | ASCII characters | | | | | | | | | |
| ← 9 bits | (O to 255) | | | | | | | | | |
| | Codes | | | | | | | | | |
| ✓ 9 bits | (256 to 512) | | | | | | | | | |
| | | | | | | | | | | |
| | | | | | | | | | | |
| 17W/ Algorithm Nightanian (1997) | aug d) | | | | | | | | | |
| LZ W AIGORITAM - DISCUSSION (CONTI | nuea) | | | | | | | | | |
| With 9 bits we can only have a maximum | n of 256 codes for | | | | | | | | | |
| strings of length 2 or above (with the f | first 256 entries for | | | | | | | | | |
| ASCII characters) | | | | | | | | | | |
| Original LZW uses dictionary with 4K e | entries, with the | | | | | | | | | |
| length of each symbol/code being 12 bi | ts | | | | | | | | | |
| · · · · · · · · · · · · · · · · · · · | ASCII characters | | | | | | | | | |
| (0 to 255 entries) | | | | | | | | | | |
| 12 bits | | | | | | | | | | |
| | Codes | | | | | | | | | |
| (256 to 1096 ontrioe) | | | | | | | | | | |
| | | | | | | | | | | |
| | | | | | | | | | | |
| \Box With 12 hits we can have a maximum of $2^{12} - 256$ codes | | | | | | | | | | |
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| Practical implementations of LZW algorithms | rithm follow the two | | | | | | | | | |
| approaches: | | | | | | | | | | |
| Flush the dictionary periodically | | | | | | | | | | |
| - no wasted codes | | | | | | | | | | |
| Grow the length of the codes as the algorithm proceeds | | | | | | | | | | |
| - First start with a length of 9 bits for the codes. | | | | | | | | | | |
| Once we run out ot codes, increase the length to 10 bits. When we run out of codes with 10 bits then we increase the code length to 11 | | | | | | | | | | |
| bits and so on. | | | | | | | | | | |
| - more efficient. | | | | | | | | | | |
| 0 ASCII | | | | | | | | | | |
| 1 Codes 256-512 0 0 ASCII | | | | | | | | | | |
| 0 1 Codes 256-511 | | | | | | | | | | |
| Out of codes 1 0 Codes 512-767 1 1 1 Codes 768-1022 | 0 0 1 Codes 256-511 | | | | | | | | | |
| | 0 1 0 Codes 512-767 | | | | | | | | | |
| | 0 1 1 Codes 768-1023 | | | | | | | | | |
| Out of codes | 1 0 0 Codes 1024-1279 | | | | | | | | | |
| | 1 0 1 Codes 1280-1535 | | | | | | | | | |
| | 1 1 0 Codes 1536-1791 | | | | | | | | | |
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