

20.1 Reviews of last lecture

In Lecture 18, we begin to study the methods for sparse recovery. Recall the basic problem is $y = W \cdot x$, where x is an $n \times 1$ high dimensional sparse data vector, W is an $m \times n$ measurement matrix with $m \ll n$, and y is an $m \times 1$ observation vector. The objective of sparse recovery can be described as: given (y, W) , find x such that $\|x\|_0 \leq S$, where $\|x\|_0$ is the number of non-zero elements in x .

We have briefly introduced three types of algorithms to perform sparse recovery:

- Recovery via l_1 -norm minimization also known as Compressed Sensing;
- A greedy algorithm also known as Orthogonal Matching Pursuit (OMP);
- Iterative methods including iterative soft thresholding and iterative hard thresholding.

We will discuss these algorithms in more details in the next few lectures. Our plan is to study l_1 -norm minimization in this lecture, OMP in the next lecture and then iterative methods.

20.2 Recovery via l_1 -norm minimization

The problem of sparse recovery via l_1 -norm minimization is often referred to as compressed sensing. A useful reference is [1] which can be found on the course website. Like we discussed earlier, the original sparse recovery problem can be described as the following Problem (20.1):

$$\text{Given } (y, W), \text{ find } x, \text{ such that } \|x\|_0 \leq S \text{ and } y = Wx. \quad (20.1)$$

Since this problem is not easy to solve in general, compressed sensing instead tries to solve an optimization problem and obtain an estimation \hat{x} of the true solution x . The optimization is defined as the following Problem (20.2):

$$\begin{aligned} & \underset{\tilde{x}}{\text{minimize}} && \|\tilde{x}\|_1 \\ & \text{subject to} && y = W\tilde{x}. \end{aligned} \quad (20.2)$$

Note in the remaining of this lecture we will always refer to x as the (true) solution to Problem (20.1), and \hat{x} as the solution to Problem (20.2).

Our hope is that we can solve the sparse recovery problem by solving compressed sensing, i.e., $\hat{x} = x$. However we quickly realize that this is not true in general, since Problem (20.2) always has a feasible solution (actually it can be written as a linear programming) while Problem (20.1) can be infeasible to solve or it can have multiple solutions. Thus in order for $\hat{x} = x$ to hold, we need to add some conditions on x and W .

Before we describe the exact conditions we need, let us see some intuitions first. Basically what we want here are the uniqueness of the sparse solution x and the hope of sparse recovery by compressed sensing.

- **Uniqueness.** If x_1 and x_2 are two solutions to Problem (20.1), then we need $x_1 = x_2$ so that we are guaranteed to have only one solution x , which is the true solution.
- **Hope of recovery.** Note that if W is such that $Wx = 0$, i.e., x is in the null space of W , then there is no hope of finding a non-zero solution using compressed sensing, since the solution of Problem (20.2) will always be $\hat{x} = 0$. Thus for sparse x not falling in the null space of W , we need some conditions that "preserves" sparse vectors, for example,

$$1 - \epsilon \leq \frac{\|W\tilde{x}\|_2}{\|\tilde{x}\|_2}, \quad \forall \tilde{x} \text{ such that } \|\tilde{x}\|_0 \leq S.$$

Now let us define **Restricted Isometry Property (RIP)** which is a sufficient condition needed for both uniqueness of sparse solution and sparse recovery by compressed sensing.

Definition 1. A matrix W satisfies (ϵ, S) -RIP if

$$1 - \epsilon \leq \frac{\|W\tilde{x}\|_2}{\|\tilde{x}\|_2} \leq 1 + \epsilon, \quad \forall \tilde{x} \text{ such that } \|\tilde{x}\|_0 \leq S.$$

Here are some remarks on the definition of RIP.

- The upper bound $1 + \epsilon$ is required in analysis, but we are not sure whether it is fundamentally required.
- Under special conditions $\epsilon = 0$, $S = n$, the above definition implies that W is an isometry on \mathbb{R}^n . This is why we use the term "restricted isometry".

We will first show uniqueness of sparse solution under RIP.

Lemma 20.1. If W satisfies $(\epsilon, 2S)$ -RIP for any $\epsilon < 1$, then $y = Wx$ cannot have two S -sparse solutions.

Proof: Suppose there exist two S -sparse solutions x_1 and x_2 . Let $\tilde{x} = x_1 - x_2$ and note that $\|W\tilde{x}\|_2 = 0$. Also note that \tilde{x} is a $2S$ -sparse vector, thus by RIP, we have $\frac{\|W\tilde{x}\|_2}{\|\tilde{x}\|_2} \geq 1 - \epsilon > 0$, which implies $\frac{0}{\|\tilde{x}\|_2} > 0$. Thus there must be $\tilde{x} = 0$, i.e., $x_1 = x_2$. \square

Note that Lemma 20.1 guarantees that under the conditions stated in the lemma, Problem (20.1) has a unique solution x , which is the true solution. Then we can state the following theorem.

Theorem 20.2. *If $\|x\|_0 \leq S$ and W satisfies $(\epsilon, 2S)$ -RIP for $\epsilon < \sqrt{2} - 1$, then $\hat{x} = x$.*

We will not prove Theorem 20.2, but instead we will prove a stronger theorem which is describe as follows:

Theorem 20.3. *If W satisfies $(\epsilon, 2S)$ -RIP with $\epsilon < \sqrt{2} - 1$, then*

$$\|x - \hat{x}\|_2 \leq \frac{2}{\sqrt{S}(1 - \rho)} \|x - x_S\|_1,$$

where $\rho = \frac{\sqrt{2}\epsilon}{1 - \epsilon}$, and x_S is the largest S entries of x .

Proof: Let \tilde{x} be any vector, and $h = \tilde{x} - x$. Define disjoint index sets $T_i \in [n]$ as follows:

- $T_0 =$ largest (in magnitude) S elements of vector x ;
- $T_1 =$ largest (in magnitude) S elements of vector $h_{T_0^c}$;
- $T_2 =$ largest (in magnitude) S elements of vector $h_{T_{0,1}^c}$;
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Note that T_0 is an index set on vector x , and by its definition we have $x_{T_0} = x_S$. But for $i \geq 1$, T_i is an index set on vector h . Under these definitions, we can state two claims which will be proved in Lemma 20.4 and Lemma 20.5 respectively.

- **Claim 1.** For any \tilde{x} such that $\|\tilde{x}\|_1 \leq \|x\|_1$, we have

$$\|h_{T_{0,1}^c}\|_2 \leq \|h_{T_0}\|_2 + \frac{2}{\sqrt{S}} \|x - x_{T_0}\|_1.$$

- **Claim 2.** For any \tilde{x} such that $W\tilde{x} = Wx$, and W satisfies $(\epsilon, 2S)$ -RIP with $\epsilon < \sqrt{2} - 1$, then

$$\|h_{T_{0,1}}\|_2 \leq \frac{\rho}{1 - \rho} \frac{1}{\sqrt{S}} \|x - x_{T_0}\|_1,$$

where $\rho = \frac{\sqrt{2}\epsilon}{1 - \epsilon}$.

Note that Claim 1 involves the l_1 -norm minimization feature, but it does not have any RIP requirements on W . Claim 2 does not involve any l_1 -norm minimization feature, but it requires RIP on W . Now if we take \tilde{x} to be the solution \hat{x} to the compressed sensing

Problem (20.2), then the conditions of Claim 1 and Claim 2 are both satisfied, and thus we have:

$$\begin{aligned}
\|x - \hat{x}\|_2 &= \|h\|_2 \leq \|h_{T_{0,1}}\|_2 + \|h_{T_{0,1}^C}\|_2 \text{ (by triangle inequality),} \\
&\leq 2\|h_{T_{0,1}}\|_2 + \frac{2}{\sqrt{S}}\|x - x_S\|_1 \text{ (by Claim 1 and the fact that } \|h_{T_0}\|_2 \leq \|h_{T_{0,1}}\|_2), \\
&\leq 2\left(\frac{\rho}{1-\rho} + 1\right)\frac{1}{\sqrt{S}}\|x - x_S\|_1 \text{ (by Claim 2),} \\
&\leq \frac{2}{\sqrt{S}(1-\rho)}\|x - x_S\|_1.
\end{aligned}$$

□

Here are the two lemmas for the two claims in the proof above.

Lemma 20.4. *With the definitions in the proof of Theorem 20.3, for any \tilde{x} such that $\|\tilde{x}\|_1 \leq \|x\|_1$, we have*

$$\|h_{T_{0,1}^C}\|_2 \leq \|h_{T_0}\|_2 + \frac{2}{\sqrt{S}}\|x - x_{T_0}\|_1.$$

Proof: Take $j > 1$, for any $i \in T_j$ and $i' \in T_{j-1}$ we have $|h_i| \leq |h_{i'}|$. Therefore, $\|h_{T_j}\|_\infty \leq \frac{\|h_{T_{j-1}}\|_1}{S}$. Thus we have:

$$\|h_{T_j}\|_2 \leq \sqrt{S}\|h_{T_j}\|_\infty \leq \frac{1}{\sqrt{S}}\|h_{T_{j-1}}\|_1.$$

Summing the above over $j = 2, 3, \dots$ and using the triangle inequality we obtain:

$$\|h_{T_{0,1}^C}\|_2 \leq \sum_{j \leq 2} \|h_{T_j}\|_2 \leq \frac{1}{\sqrt{S}}\|h_{T_0^C}\|_1. \quad (20.3)$$

Since $\|\tilde{x}\|_1 \leq \|x\|_1$, using triangle inequality we have:

$$\|x\|_1 \geq \|x + h\|_1 = \sum_{i \in T_0} |x_i + h_i| + \sum_{i \in T_0^C} |x_i + h_i| \geq \|x_{T_0}\|_1 - \|h_{T_0}\|_1 + \|h_{T_0^C}\|_1 - \|x_{T_0^C}\|_1, \quad (20.4)$$

and since $\|x_{T_0^C}\|_1 = \|x - x_S\|_1 = \|x\|_1 - \|x_{T_0}\|_1$, Equation (20.4) can be further written as:

$$\|h_{T_0^C}\|_1 \leq \|h_{T_0}\|_1 + 2\|x_{T_0^C}\|_1. \quad (20.5)$$

Combining Equation (20.3) and (20.5), we obtain:

$$\|h_{T_{0,1}^C}\|_2 \leq \frac{1}{\sqrt{S}}(\|h_{T_0}\|_1 + 2\|x_{T_0^C}\|_1) \leq \|h_{T_0}\|_2 + \frac{2}{\sqrt{S}}\|x_{T_0^C}\|_1.$$

□

Lemma 20.5. *With the definitions in the proof of Theorem 20.3, for any \tilde{x} such that $W\tilde{x} = Wx$, and W satisfies $(\epsilon, 2S)$ -RIP with $\epsilon < \sqrt{2} - 1$, then*

$$\|h_{T_{0,1}}\|_2 \leq \frac{\rho}{1 - \rho} \frac{1}{\sqrt{S}} \|x - x_{T_0}\|_1,$$

where $\rho = \frac{\sqrt{2}\epsilon}{1 - \epsilon}$.

Proof: Since vector $h_{T_{0,1}}$ is $2S$ -sparse, we can use the RIP condition to get:

$$(1 - \epsilon)\|h_{T_{0,1}}\|_2^2 \leq \|Wh_{T_{0,1}}\|_2^2. \quad (20.6)$$

Note that $Wh_{T_{0,1}} = Wh - \sum_{j \geq 2} Wh_{T_j} = -\sum_{j \geq 2} Wh_{T_j}$, thus we have:

$$\|Wh_{T_{0,1}}\|_2^2 = -\sum_{j \geq 2} \langle Wh_{T_{0,1}}, Wh_{T_j} \rangle = -\sum_{j \geq 2} \langle Wh_{T_0} + Wh_{T_1}, Wh_{T_j} \rangle. \quad (20.7)$$

Note that for all i and j , $i \neq j$, $\frac{h_{T_i}}{\|h_{T_i}\|_2} + \frac{h_{T_j}}{\|h_{T_j}\|_2}$ and $\frac{h_{T_i}}{\|h_{T_i}\|_2} - \frac{h_{T_j}}{\|h_{T_j}\|_2}$ are $2S$ -sparse, and then by RIP we have:

$$\left\| \frac{Wh_{T_i}}{\|h_{T_i}\|_2} + \frac{Wh_{T_j}}{\|h_{T_j}\|_2} \right\|_2^2 \leq (1 + \epsilon) \left(\left\| \frac{h_{T_i}}{\|h_{T_i}\|_2} \right\|_2^2 + \left\| \frac{h_{T_j}}{\|h_{T_j}\|_2} \right\|_2^2 \right) = 2(1 + \epsilon),$$

and

$$-\left\| \frac{Wh_{T_i}}{\|h_{T_i}\|_2} - \frac{Wh_{T_j}}{\|h_{T_j}\|_2} \right\|_2^2 \leq -(1 - \epsilon) \left(\left\| \frac{h_{T_i}}{\|h_{T_i}\|_2} \right\|_2^2 + \left\| \frac{h_{T_j}}{\|h_{T_j}\|_2} \right\|_2^2 \right) = -2(1 - \epsilon).$$

Thus we have:

$$\left| \left\langle \frac{Wh_{T_i}}{\|h_{T_i}\|_2}, \frac{Wh_{T_j}}{\|h_{T_j}\|_2} \right\rangle \right| = \left| \frac{\left\| \frac{Wh_{T_i}}{\|h_{T_i}\|_2} + \frac{Wh_{T_j}}{\|h_{T_j}\|_2} \right\|_2^2 - \left\| \frac{Wh_{T_i}}{\|h_{T_i}\|_2} - \frac{Wh_{T_j}}{\|h_{T_j}\|_2} \right\|_2^2}{4} \right| \leq \epsilon,$$

which implies:

$$|\langle Wh_{T_i}, Wh_{T_j} \rangle| \leq \epsilon \|h_{T_i}\|_2 \|h_{T_j}\|_2.$$

Take the above equation into Equation (20.7), we obtain:

$$\|Wh_{T_{0,1}}\|_2^2 \leq \epsilon (\|h_{T_0}\|_2 + \|h_{T_1}\|_2) \sum_{j \geq 2} \|h_{T_j}\|_2 \leq \sqrt{2}\epsilon \|h_{T_{0,1}}\|_2 \sum_{j \geq 2} \|h_{T_j}\|_2. \quad (20.8)$$

Combining Equation (20.3) (20.6) and (20.8), we obtain:

$$(1 - \epsilon)\|h_{T_{0,1}}\|_2^2 \leq \sqrt{2}\epsilon \|h_{T_{0,1}}\|_2 \frac{1}{\sqrt{S}} \|h_{T_0^c}\|_1.$$

Using Equation (20.5) and rearranging, we get:

$$\|h_{T_0,1}\|_2 \leq \frac{\rho}{\sqrt{S}}(\|h_{T_0}\|_1 + 2\|x_{T_0^c}\|_1) \leq \rho\|h_{T_0}\|_2 + \frac{2\rho}{\sqrt{S}}\|x_{T_0^c}\|_1,$$

but since $\|h_{T_0}\|_2 \leq \|h_{T_0,1}\|_2$, we have:

$$\|h_{T_0,1}\|_2 \leq \frac{\rho}{1-\rho} \frac{1}{\sqrt{S}} \|x - x_{T_0}\|_1.$$

□

Bibliography

- [1] Shai Shalev-Shwartz. Compressed sensing: Basic results and self contained proofs. in manuscript, 2009.