## Exam 1

## Date: October 3, 2013

UT EID:		
Printed Name:	Last,	First
Your signature is on this exam:	your promise that you have not cheated and w	ill not cheat on this exam, nor will you help others to cheat
Signature:		

## **Instructions:**

- Closed book and closed notes. No books, no papers, no data sheets (other than the last two pages of this Exam)
- No devices other than pencil, pen, eraser (no calculators, no electronic devices), please turn cell phones off.
- Please be sure that your answers to all questions (and all supporting work that is required) are contained in the space (boxes) provided. *Anything outside the boxes will be ignored in grading*.
- You have 75 minutes, so allocate your time accordingly.
- For all questions, unless otherwise stated, find the most efficient (time, resources) solution.
- Unless otherwise stated, make all I/O accesses friendly.
- Please read the entire exam before starting.

Problem 1	10	
Problem 2	10	
Problem 3	10	
Problem 4	10	
Problem 5	10	
Problem 6	15	
Problem 7	25	
Problem 8	10	
Total	100	

**Part a)** A last in first out data storage system on the computer used to remember data temporarily.

**Part b)** Software is added to the system for the purpose of debugging, and this software has a large and significant effect on the system.

Part c) This C operator is used to perform a right shift.

**Part d)** The name given to describe  $1,000 (10^3)$  bytes.

**Part e)** A type of logic where the voltage representing false is less than the voltage representing true.

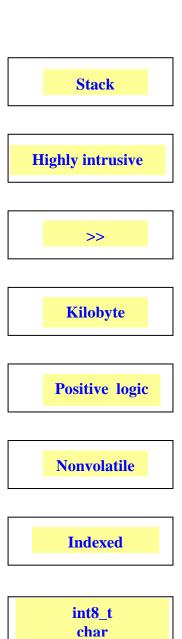
**Part f)** A property of ROM such that data is not lost if power is removed and then restored.

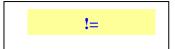
**Part g)** This addressing mode is always used to access memory, shown here as the source operand of this instruction: LDR R1, [R0]?

**Part h**) This declaration is used to create a variable in C that has a precision of 8 bits and can take negative values.

**Part i)** This C operator is used in if-then while-loop and do-while-loops for checking to see if two numbers are not equal.

**Part j)** A drawing that describes the sequence of operations of software, defining what and when software actions will occur.







Question 2 (10 points) Consider the following 8-bit subtraction (assume registers are 8 bits wide) Load 0xC0 into R1

```
Load 0x1F into R2
Subtract R3 = R1-R2
```

**a.** What will be the 8-bit result in Register R3 (in hex)?

```
<mark>0xC0 - 0x1F = 0xA1</mark>
```

**b.** What is 8-bit result in Register R3 (as an unsigned decimal)?

```
10*16+1=161
```

c. What is 8-bit result in Register R3 (as a signed decimal)?

```
-128+32+1=-95
```

**d.** What will be the value of the carry (C) bit?

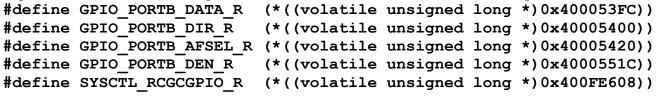
```
192 - 31 = 161
the unsigned answer is correct, C=1
```

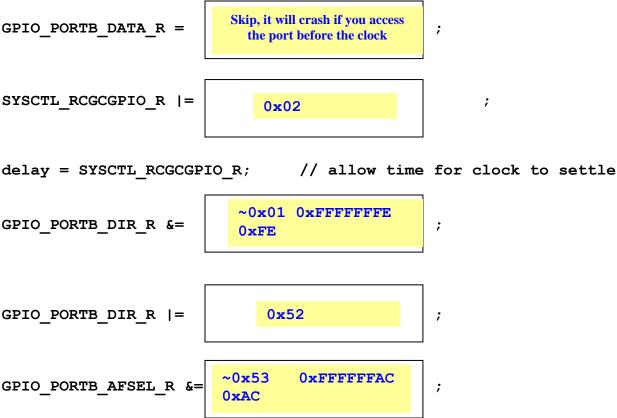
e. What will be the value of the overflow (V) bit?

```
-64 - 31 = -95
the signed answer is correct, V=0
```

GPIO PORTB DEN R |=

(10) Question 3. You will fill in the blanks of this C code that initializes Port B. Make pins PB6, PB4, PB1 outputs. Make the pin PB0 an input. To get full credit, this code must be friendly. Partial credit can be obtained by writing code that works, but is not friendly. Leave the box empty if it is not required to be executed at that spot in the initialization. You will use the following definitions:

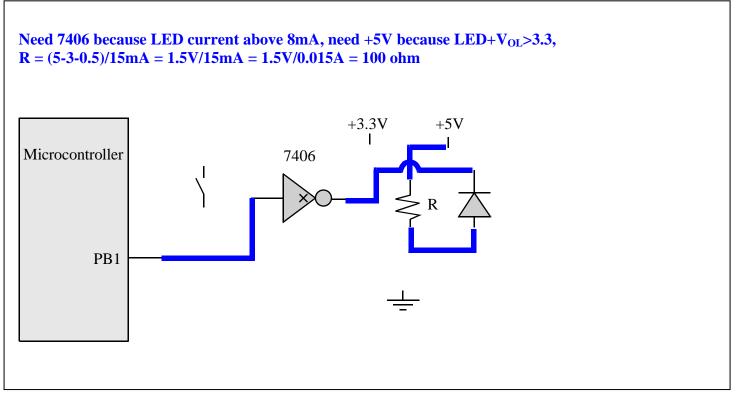




0x53

(10) Question 4. Interface the LED to PB1 such that if PB1 is high, the LED is on, and if PB1 is low the LED is off. The desired LED operating point is 3.0V at 15 mA. The  $V_{OH}$  of the microcontroller is 3.1 V. The  $V_{OL}$  of the microcontroller is 0.3 V. The maximum current that the microcontroller can source or sink is 8 mA. The  $V_{OL}$  of the 7406 is 0.5 V. The maximum current that the 7406 can sink is 40 mA. Your bag of parts includes the switch, the 7406, the LED, and a resistor (you specify the resistor value). Pick the fewest components to use. You will not need them all. You may also use 3.3V, 5V power, and/or ground. Show the equations used to calculate the resistor value.

;



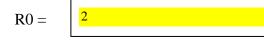
(10) Question 5. Write an assembly subroutine, called Calc, that calculates Output = Input/8 -5. The *Input* and *Output* parameters are 8-bit signed numbers located in global RAM. You may use Registers R0-R3, or R12 as scratch registers without saving and restoring them. Full credit will be given to the fastest solution. Don't worry about the code to define *Input* and *Output*, just the subroutine.

```
; Calculate Output = Input/8 -5
          R0 (-128 to +127)
; Input:
; Output: R0 (-21 to +10), can't underflow or overflow
Calc
      LDR
            R1,=Input
      LDRSB R0, [R1]
      ASR
            R0,R0,#3
                       ; Input/8 (signed)
                       ; Input/8-5
      SUB
            R0,R0,#5
      LDR
            R2,=Output
      STRB
            R0,[R2]
      BX
            LR
                       ; return
```

(15) Question 6. Answer the following questions with reference to the C and assembly code below. You may assume that all linkages have been done to be able to call the assembly code from C. <u>Hint</u>: Recall AAPCS

; C code calling assembly	; As	sembly	y code
<pre>int32_t Param1;</pre>	Bs	rn 0	
int32_t Param2;	Res	RN 1	
int32_t Output;	Ex	RN 4	
	Prod	RN 5	
<pre>int main() {</pre>	Sub	PUSH	{R4,R5,LR}
Param1 = 2; Param2 = 7;		MOV	Ex,#0
Output = Sub(Param1,Param2);		MOV	Prod,Bs
}	More	CMP	Prod,Res
		BGT	Done
		MUL	Prod, Prod, Bs
		ADD	Ex,Ex,#1
		в	More
	Done	MOV	R0,Ex
		POP	{R4,R5,LR}
		вх	LR

(2) Part a) What is the numerical value in register R0 at the start of the assembly subroutine Sub?



(2) Part b) What is the numerical value in register R1 at the start of the assembly subroutine Sub?  $R1 = \begin{bmatrix} 7 \\ \hline 7 \\ \hline \end{bmatrix}$ 

(4) Part c) What is the numerical value of the C variable Output after the assignment statement,Output = Sub(Param1, Param2); is executed?

output =

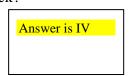
(2) Part d) Why did the subroutine **Sub**, save the registers R4 and R5 on the stack?

- I. The input parameters are on the stack.
- II. The output parameter is returned on the stack
- III. In order to save the return address
- IV. Follows AAPCS convention
- V. None of the above.

(5) Part e) Which of the following statements describes what Sub does accurately?

- I. Sub returns the product of the two inputs using successive addition
- II. Sub returns the largest power to which the first input can be raised and still have it less than or equal to the second
- III. Sub returns the exponent of the first input raised to the second input
- IV. Sub returns the smallest power to which the first input needs to be raised so that it is greater than or equal to the second
- V. Sub returns the power to which the first input has to be raised to be equal to the second

Answer is II



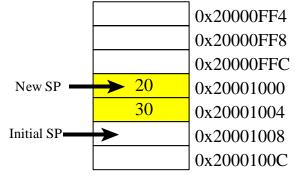
(20) Question 7. You are asked to develop the software for a control panel of a home automation system. *You can write software in either assembly or C*. Make sure that all of your software is friendly and follows the AAPCS. You may assume the hardware is already connected, and Port B is already initialized so PB5 is an output and PB4–0 are inputs. Please use the port definition GPIO\_PORTB\_DATA\_R to access Port B. You are not allowed to use bit-specific port addressing. The system has five door/window switches (sensors) connected to pins PB4, PB3, PB2, PB1, PB0. Door/window signals are high if OK, and low if there is danger. There is an LED connected to pin PB5, which signifies an alarm. The LED interface is negative logic. Write the main program of the control panel that continuously checks the sensors and turns the warning LED connected if, and only if, two or more door/window switches indicate there is danger.

```
void main(void) { uint32 t count,mask;
  while(1){
    count = 0;
    for (mask = 0x01; mask < 0x20; mask=mask<<1) {
      if((GPIO PORTB DATA R&mask)==0) { // danger if low
        count++;
      }
    }
    if(count > 1){
      GPIO PORTE DATA R &= \sim 0x20; // negative logic LED on
    } else{
      GPIO PORTB DATA R |= 0x20; // negative logic LED off
    }
  }
}
void main(void) { uint32 t count;
  while(1){
    count = 0;
    if ((GPIO PORTB DATA R&0x01) == 0) count++; // PB0 not pressed
    if((GPIO PORTB DATA R&0x02)==0) count++; // PB1 not pressed
    if((GPIO PORTB DATA R&0x04)==0) count++; // PB2 not pressed
    if((GPIO PORTB DATA R&0x08)==0) count++; // PB3 not pressed
    if((GPIO PORTB DATA R&0x10)==0) count++; // PB4 not pressed
    if(count > 1){
      GPIO PORTB DATA R &= \sim 0 \times 20; // negative logic LED on
    } else{
      GPIO PORTB DATA R |= 0 \times 20; // negative logic LED off
    }
  }
}
```

(10) Question 8. Show the contents of the stack after the two marked points in the execution of the following code. Assume R0=0, R1=10, R2=20, R3=30, R4=40, R5=50, and R6=60. The initial stack pointer is 0x20001008.

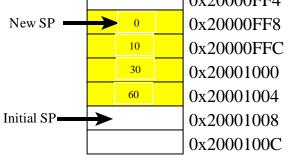
PUSH	{R2,R3}
ADD	R4,R1,R0 ; < A
POP	{R5,R6}
ADD	R5,R4
ADD	R6,R5
PUSH	{R0,R4-R6};
	< В

a) (4 points) The contents of the stack (SP and contents) after execution point A:



b) (6 points) The contents of the stack (SP and contents) after execution point B:

PUSH {R2,R3} ADD R4,R1,R0 ; R4=10 POP {R5,R6} ; R5=20, R6=30 ADD R5,R4 ; R5=10+20=30 ADD R6,R5 ; R6=30+30=60 PUSH {R0,R4-R6}; push 0,10,30,60 <---- B 0x20000FF4



```
Memory access instructions
   LDR
          Rd, [Rn]
                        ; load 32-bit number at [Rn] to Rd
   LDR
          Rd, [Rn, #off] ; load 32-bit number at [Rn+off] to Rd
          Rd, =value ; set Rd equal to any 32-bit value (PC rel)
   LDR
                        ; load unsigned 16-bit at [Rn] to Rd
   LDRH
          Rd, [Rn]
          Rd, [Rn,#off] ; load unsigned 16-bit at [Rn+off] to Rd
  LDRH
   LDRSH Rd, [Rn]
                         ; load signed 16-bit at [Rn] to Rd
   LDRSH Rd, [Rn, #off] ; load signed 16-bit at [Rn+off] to Rd
                        ; load unsigned 8-bit at [Rn] to Rd
   LDRB
          Rd, [Rn]
   LDRB
          Rd, [Rn,#off] ; load unsigned 8-bit at [Rn+off] to Rd
   LDRSB Rd, [Rn] ; load signed 8-bit at [Rn] to Rd
  LDRSB Rd, [Rn, #off] ; load signed 8-bit at [Rn+off] to Rd
          Rt, [Rn] ; store 32-bit Rt to [Rn]
Rt, [Rn,#off] ; store 32-bit Rt to [Rn+off]
   STR
   STR
   STRH
         Rt, [Rn] ; store least sig. 16-bit Rt to [Rn]
   STRH
          Rt, [Rn,#off] ; store least sig. 16-bit Rt to [Rn+off]
   STRB
          Rt, [Rn] ; store least sig. 8-bit Rt to [Rn]
          Rt, [Rn,#off] ; store least sig. 8-bit Rt to [Rn+off]
   STRB
                 ; push 32-bit Rt onto stack
   PUSH
         {Rt}
  POP{Rd}; pop 32-bit number from stack into RdADRRd, label; set Rd equal to the address at labelMOV{S} Rd, <op2>; set Rd equal to op2MOVRd, #im16; set Rd equal to im16, im16 is 0 to 65535MVN{S} Rd, <op2>; set Rd equal to -op2
Branch instructions
        label ; branch to label
  в
                                      Always
   BEQ label ; branch if Z == 1
                                      Equal
  BNE label ; branch if Z == 0
                                      Not equal
  BCS label ; branch if C == 1
                                      Higher or same, unsigned \geq
  BHS label ; branch if C == 1
                                     Higher or same, unsigned \geq
  BCC label ; branch if C == 0
                                      Lower, unsigned <
  BLO label ; branch if C == 0
                                      Lower, unsigned <
  BMI label ; branch if N == 1
                                      Negative
  BPL label ; branch if N == 0
                                      Positive or zero
  BVS label ; branch if V == 1
                                      Overflow
  BVC label ; branch if V == 0
                                      No overflow
  BHI label ; branch if C==1 and Z==0 Higher, unsigned >
  BLS label ; branch if C==0 or Z==1 Lower or same, unsigned \leq
  BGE label ; branch if N == V
                                      Greater than or equal, signed \geq
  BLT label ; branch if N != V
                                      Less than, signed <
  BGT label ; branch if Z==0 and N==V Greater than, signed >
  BLE label ; branch if Z==1 or N!=V Less than or equal, signed \leq
               ; branch indirect to location specified by Rm
   BX
        Rm
   BL
        label
              ; branch to subroutine at label
   BLX Rm ; branch to subroutine indirect specified by Rm
Interrupt instructions
   CPSIE I
                          ; enable interrupts (I=0)
   CPSID I
                          ; disable interrupts (I=1)
Logical instructions
   AND\{S\} {Rd,} Rn, <op2> ; Rd=Rn&op2
                                          (op2 is 32 bits)
   ORR{S} {Rd,} Rn, <op2> ; Rd=Rn|op2
EOR{S} {Rd,} Rn, <op2> ; Rd=Rn^op2
                                          (op2 is 32 bits)
                                          (op2 is 32 bits)
   BIC{S} {Rd,} Rn, <op2> ; Rd=Rn&(~op2) (op2 is 32 bits)
   ORN{S} {Rd,} Rn, <op2> ; Rd=Rn|(~op2) (op2 is 32 bits)
   LSR{S} Rd, Rm, Rs ; logical shift right Rd=Rm>>Rs (unsigned)
```

```
LSR{S} Rd, Rm, #n
                            ; logical shift right Rd=Rm>>n
                                                                 (unsigned)
   ASR{S} Rd, Rm, Rs
                            ; arithmetic shift right Rd=Rm>>Rs (signed)
                           ; arithmetic shift right Rd=Rm>>n (signed)
   ASR{S} Rd, Rm, #n
   LSL{S} Rd, Rm, Rs
                           ; shift left Rd=Rm<<Rs (signed, unsigned)
   LSL{S} Rd, Rm, #n
                           ; shift left Rd=Rm<<n (signed, unsigned)</pre>
Arithmetic instructions
   ADD{S} {Rd,} Rn, \langle op2 \rangle; Rd = Rn + op2
   ADD{S} {Rd}, Rn, \#im12; Rd = Rn + im12, im12 is 0 to 4095
   SUB{S} {Rd,} Rn, \langle op2 \rangle; Rd = Rn - op2
   SUB{S} {Rd,} Rn, \#im12; Rd = Rn - im12, im12 is 0 to 4095
   RSB{S} {Rd}, Rn, <op2>; Rd = op2 - Rn
   RSB{S} {Rd_{i}} Rn_{i} \#im12 ; Rd = im12 - Rn
                        ; Rn – op2
   CMP
          Rn, <op2>
                                              sets the NZVC bits
   CMN
          Rn, <op2>
                            ; Rn - (-op2)
                                              sets the NZVC bits
                          ; Rd = Rn * Rm
   MUL{S} {Rd,} Rn, Rm
                                                   signed or unsigned
          Rd, Rn, Rm, Ra; Rd = Ra + Rn*Rm
                                                   signed or unsigned
   MLA
   MLS
          Rd, Rn, Rm, Ra ; Rd = Ra - Rn*Rm
                                                   signed or unsigned
   UDIV
           {Rd,} Rn, Rm
                            ; Rd = Rn/Rm
                                                   unsigned
   SDTV
           {Rd,} Rn, Rm
                            ; Rd = Rn/Rm
                                                   signed
Notes Ra Rd Rm Rn Rt represent 32-bit registers
              any 32-bit value: signed, unsigned, or address
     value
              if S is present, instruction will set condition codes
     {S}
     #im12
              any value from 0 to 4095
     #im16
              any value from 0 to 65535
              if Rd is present Rd is destination, otherwise Rn
     \{Rd_{\ell}\}
     #n
              any value from 0 to 31
     #off
              any value from -255 to 4095
              any address within the ROM of the microcontroller
     label
              the value generated by <op2>
     op2
Examples of flexible operand <op2> creating the 32-bit number. E.g., Rd = Rn+op2
   ADD Rd, Rn, Rm
                            ; op2 = Rm
   ADD Rd, Rn, Rm, LSL #n ; op2 = Rm<<n Rm is signed, unsigned
   ADD Rd, Rn, Rm, LSR #n ; op2 = Rm>>n Rm is unsigned
   ADD Rd, Rn, Rm, ASR #n ; op2 = Rm>>n Rm is signed
   ADD Rd, Rn, #constant ; op2 = constant, where X and Y are hexadecimal digits:
             •
                produced by shifting an 8-bit unsigned value left by any number of bits
                in the form 0x00XY00XY
             •
                in the form 0xXY00XY00
                in the form 0xXYXYXYX
                  R0
                                                                             0x0000.0000
                  R1
                                                                 256k Flash
                  R2
                            Condition code bits
                                                                   ROM
                  R3
                                                                             0x0003.FFFF
                            N negative
                  R4
    General
                  R5
                            Z zero
                                                                             0x2000.0000
                  R6
                                                                  32k RAM
   purpose -
                            V signed overflow
   registers
                  R7
                                                                             0x2000.7FFF
                            C carry or
                  R8
                  R9
                              unsigned overflow
                                                                             0x400<u>0</u>.0000
                 R10
                                                                  I/O ports
                  R11
                  R12
                                                                             0x400F.FFFF
    Stack pointer
              R13 (MSP)
    Link register
                R14 (LR)
                                                                             0xE000.0000
  Program counter R15 (PC)
                                                                 Internal I/O
                                                                    PPB
                                                                             0xE004.1FFF
```