Exam 1

Date: February 20, 2014

UT EID:		
Printed Name:	Last,	First
Your signature is on this exam:	your promise that you have not cheated and	will not cheat on this exam, nor will you help others to cheat
Signature:		

Instructions:

- Closed book and closed notes. No books, no papers, no data sheets (other than the last two pages of this Exam)
- No devices other than pencil, pen, eraser (no calculators, no electronic devices), please turn cell phones off.
- Please be sure that your answers to all questions (and all supporting work that is required) are contained in the space (boxes) provided. *Anything outside the boxes/blanks will be ignored in grading*. You may use the back of the sheets for scratch work.
- You have 75 minutes, so allocate your time accordingly.
- For all questions, unless otherwise stated, find the most efficient (time, resources) solution.
- Unless otherwise stated, make all I/O accesses friendly.
- Please read the entire exam before starting.

Problem 1	10	
Problem 2	10	
Problem 3	10	
Problem 4	10	
Problem 5	10	
Problem 6	15	
Problem 7	25	
Problem 8	10	
Total	100	

(10) Question 1. Place your answers in the boxes.

Part a) The total addressable memory on the ARM Cortex M processor is these many bytes. Give the total possible, not the actual number on TM4C123.

Part b) This C operator is used to perform a bit-wise NOT operation is.

Part c) We access device registers just like we access memory. The term used for this kind of I/O is.

Part d) In conditional C expressions, a Zero value is interpreted as False. What is interpreted a True?

Part e) What LED parameter (Voltage or Current) determines whether we need a 7406 driver.

Part f) A big-endian machine will interpret a byte 0x28 as having a value of (2*16+8) = 40, whereas a little-endian machine will interpret 0x28 as (8*16+2) = 130. *True or False*.

Part g) What is the term "addressing mode" associated with, *Instructions* or *Operands*?

Part h) This data-type is the most appropriate one to create a variable in C that can take values in the range -200 to +200.

Part i) Which three registers or the ARM Cortex processor are initialized on Reset?

Part j) Give an example of a non-intrusive debugging tool.









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(10) Question 2

Consider the following arithmetic operation. Assume that all registers are 8-bit.

 $R0 \leftarrow 100$ R1 \leftarrow 227 R2 \leftarrow R0 - R1

What are the numbers in registers R1 and R2 in unsigned decimal and signed decimal?

[R1] = [R1] =	 (unsigned decimal) (signed decimal)
[R2] = [R2] =	 (unsigned decimal) (signed decimal)

What are the values of the condition code bits: N, Z, V, C?

N = _____ Z = _____ V = _____ C = _____

(10) Question 3

A programmer wants to make make pins PB1, PB4, PB7 outputs and make pin PB0 an input. So he writes the below code in sequence as shown. He claims it does not work. Your job as an expert is to identify the mistake(s). First, start by commenting what is the purpose of each statement in the code. Second, after each statement, please write either OK or explain what is wrong and provide the necessary corrections to make things work as expected. You are free to add, remove or modify the sequence, as well as the code. Assume that you have access to the following correct definitions:

```
#define GPIO_PORTB_DATA_R (*((volatile unsigned long *)0x400053FC))
#define GPIO_PORTB_DIR_R (*((volatile unsigned long *)0x40005400))
#define GPIO_PORTB_AFSEL_R (*((volatile unsigned long *)0x40005420))
#define GPIO_PORTB_DEN_R (*((volatile unsigned long *)0x4000551C))
#define SYSCTL_RCGCGPIO_R (*((volatile unsigned long *)0x400FE608))
```

SYSCTL_RCGCGPIO_R = 0×02 ;

 $GPIO_PORTB_DATA_R = 0 \times 00;$

GPIO_PORTB_DIR_R |= 0x92;

GPIO_PORTB_DIR_R &= 0x01;

GPIO_PORTB_AFSEL_R &= ~0x93;

(10) Question 4.

(a) Interface a switch using a 10 k Ω resistor to port PA5 using positive logic. You may assume PA5 has been configured as an input port. Also assume that no current can flow into and out of the port pin and the switch is ideal. Find the current through the switch and the voltage across the resistor. Complete the table below the figure.



(b) Interface an LED through a resistor to port PA4 using negative logic. You may assume PA4 has been configured as an output port. The operating point of this LED is 1.5V at 1.8mA. The V_{OL} and V_{OH} of the TM4C123 is 0.3V and 3.3V resp., and the maximum current that PA4 can source or sink is 8mA. Find the value of the resistor R that needs to be connected, and show the circuit diagram. What is the current through the LED and the voltage across the resistor? Complete the table below the figure.





PA4 output	Current through LED	Voltage across the resistor
Logic level 1		
Logic level 0		

(10) Question 5. Write an *assembly* subroutine, called SwapLT, that swaps the contents of two global variables *ying* and *yang* only if *ying* is less than *yang*. Assume that variables are 16-bit signed numbers located in global RAM. You may use Registers R0-R3, or R12 as scratch registers without saving and restoring them. Don't worry about how the variables are initialized.

```
AREA DATA, ALIGN=2

ying SPACE 2

yang SPACE 2

AREA |.text|, CODE, READONLY, ALIGN=2
```

(15) Question 6. You are given the assembly Subroutine Max3 below that takes three inputs and returns the maximum of the three inputs. one two three are input parameters, and max is the return parameter. There are no bugs in this subroutine, but it is not AAPCS compliant. Part a) Make changes to anything inside the box so Max3 becomes AAPCS compliant.

RN 1 one RN 2 two 3 three RN RN 5 max EXPORT Max3 Max3 PUSH {R4} CMP one, two BHI C13 CMP two,three BHI Mx2 В Mx3 C13 CMP one, three BLO Mx3 MOV max, one В Done MOV max, two Mx2 В Done Mx3 MOV max, three Done POP $\{R4\}$ BX LR

Assume these global variables uint32_t result,w,x,y,z;

Part b) Write one line of C code that calls the **Max3** assembly subroutine passing the values 7000, 1134, and 4556 storing the result in global variable **result**.

Part c) Write one line of C code that calls the **Max3** assembly subroutine passing the contents of global variables **x**, **y**, and **z** storing the result in global variable **w**.

(25) Question 7. You are asked to write a software module that controls the child-lock feature on car door. An indicator light (connected to PortA pin 1 on the microcontroller) on the dashboard shows whether the child-lock feature is engaged or not. A switch (connected to PortA pin 0) controls whether the feature is enabled or disabled. There is some external hardware in the form of a weight sensor that sends an 8-bit input value on PortB indicating the weight of the person/object in the seat. If the child-lock feature is enabled by turning the switch (PA0) to ON, you have to read the weight from PortB and check it to see if the indicator light must be turned on. The indicator light (on PA1) must be turned on when the weight (8-bit value on PortB) is between 10 (GPIO_PORTB_DATA_R=0x0A) and 35 (GPIO_PORTB_DATA_R=0x23) and the child-lock feature is enabled (PA0 is 1). Otherwise, it must be off.

You may assume the hardware is already connected, and all initializations except setting/clearing the Direction register for Port B, are done. You have access to registers **GPIO_PORTB_DIR_R**, **GPIO_PORTA_DATA_R**, and **GPIO_PORTB_DATA_R** to complete your code. Write **assembly** code that manipulates the direction registers first and then continuously checks the switch and weight and updates the LED state accordingly.

Suggestion: It might help if you visualize the solution using a flowchart

; Declarations already in place for GPIO registers Start ...; Code for GPIO initialization you are not responsible for is here ; *** Code you are responsible for follows **** ; Port B Direction register initialization ; Logic to check inputs and produce appropriate outputs follows Loop

B Loop ALIGN END

Initial SP

(10) Question 8. Show the contents of the stack after the two marked points in the execution of the following code. The initial stack pointer is 0×20001008 . 0×00002000 MOV B0 #3 Sub

070002000	HOV KO,	#J		Sub	
0x00002002	MUL R4,	R0,R0		PUSH	$\{LR, R4, R5, R6\}$
0x00002004	ADD R5,	R4,#1		MUL	R4,R4,R0 ; < A
0x00002006	ADD R6,	R4,#2		POP	{R4,R5,R6,PC}
0x00002008	BL Sub				
0x0000200A	ADD R4	,R5,R6			
		;	< В		
a) (4 points)	Give the state of	of the stack	k (SP and contents) after exe	cution poin	nt A:
		0x2000	0FF4		
		0x2000	0FF8		
		0x2000	0FFC		
		0x2000	1000		
		0x2000	1004		

0x20001008

0x2000100C

b) (6 points) Give the state of the stack (SP and contents) after execution point B and the value stored in R4:



R4 =	
1	

Memory access instructions LDR Rd, [Rn] ; load 32-bit number at [Rn] to Rd LDR Rd, [Rn, #off] ; load 32-bit number at [Rn+off] to Rd Rd, =value ; set Rd equal to any 32-bit value (PC rel) LDR ; load unsigned 16-bit at [Rn] to Rd LDRH Rd, [Rn] Rd, [Rn,#off] ; load unsigned 16-bit at [Rn+off] to Rd LDRH LDRSH Rd, [Rn] ; load signed 16-bit at [Rn] to Rd LDRSH Rd, [Rn,#off] ; load signed 16-bit at [Rn+off] to Rd LDRB Rd, [Rn] ; load unsigned 8-bit at [Rn] to Rd LDRB Rd, [Rn,#off] ; load unsigned 8-bit at [Rn+off] to Rd LDRSB Rd, [Rn] ; load signed 8-bit at [Rn] to Rd LDRSB Rd, [Rn, #off] ; load signed 8-bit at [Rn+off] to Rd Rt, [Rn] ; store 32-bit Rt to [Rn] Rt, [Rn,#off] ; store 32-bit Rt to [Rn+off] STR STR STRH Rt, [Rn] ; store least sig. 16-bit Rt to [Rn] STRH Rt, [Rn,#off] ; store least sig. 16-bit Rt to [Rn+off] STRB Rt, [Rn] ; store least sig. 8-bit Rt to [Rn] STRB Rt, [Rn, #off] ; store least sig. 8-bit Rt to [Rn+off] ; push 32-bit Rt onto stack PUSH {Rt} POP{Rd}; pop 32-bit number from stack into RdADRRd, label; set Rd equal to the address at labelMOV{S} Rd, <op2>; set Rd equal to op2MOVRd, #im16; set Rd equal to im16, im16 is 0 to 65535MVN{S} Rd, <op2>; set Rd equal to -op2 Branch instructions label ; branch to label в Always BEQ label ; branch if Z == 1Equal BNE label ; branch if Z == 0Not equal BCS label ; branch if C == 1 Higher or same, unsigned \geq BHS label ; branch if C == 1 Higher or same, unsigned \geq BCC label ; branch if C == 0 Lower, unsigned < BLO label ; branch if C == 0 Lower, unsigned < BMI label ; branch if N == 1 Negative BPL label ; branch if N == 0 Positive or zero BVS label ; branch if V == 1 Overflow BVC label ; branch if V == 0No overflow BHI label ; branch if C==1 and Z==0 Higher, unsigned > BLS label ; branch if C==0 or Z==1 Lower or same, unsigned \leq BGE label ; branch if N == VGreater than or equal, signed \geq BLT label ; branch if N != V Less than, signed < BGT label ; branch if Z==0 and N==V Greater than, signed > BLE label ; branch if Z==1 or N!=V Less than or equal, signed \leq ; branch indirect to location specified by Rm BX Rm BL label ; branch to subroutine at label BLX Rm ; branch to subroutine indirect specified by Rm Interrupt instructions CPSIE I ; enable interrupts (I=0) CPSID I ; disable interrupts (I=1) Logical instructions $AND{S} {Rd}, Rn, <op2>; Rd=Rn&op2$ (op2 is 32 bits) ORR{S} {Rd,} Rn, <op2> ; Rd=Rn|op2 EOR{S} {Rd,} Rn, <op2> ; Rd=Rn^op2 (op2 is 32 bits) (op2 is 32 bits) BIC{S} {Rd,} Rn, <op2> ; Rd=Rn&(~op2) (op2 is 32 bits) ORN{S} {Rd,} Rn, <op2> ; Rd=Rn|(~op2) (op2 is 32 bits) LSR{S} Rd, Rm, Rs ; logical shift right Rd=Rm>>Rs (unsigned)

LSR{S} Rd, Rm, #n ; logical shift right Rd=Rm>>n (unsigned) ASR{S} Rd, Rm, Rs ; arithmetic shift right Rd=Rm>>Rs (signed) ; arithmetic shift right Rd=Rm>>n (signed) ASR{S} Rd, Rm, #n LSL{S} Rd, Rm, Rs ; shift left Rd=Rm<<Rs (signed, unsigned)</pre> LSL{S} Rd, Rm, #n ; shift left Rd=Rm<<n (signed, unsigned)</pre> **Arithmetic instructions** ADD{S} {Rd,} Rn, $\langle op2 \rangle$; Rd = Rn + op2 $ADD{S} {Rd}, Rn, \#im12$; Rd = Rn + im12, im12 is 0 to 4095 SUB{S} {Rd,} Rn, $\langle op2 \rangle$; Rd = Rn - op2 $SUB{S} {Rd}, Rn, \#im12$; Rd = Rn - im12, im12 is 0 to 4095 $RSB{S} {Rd}, Rn, <op2>; Rd = op2 - Rn$ $RSB{S} {Rd_{,}} Rn_{,} \#im12 ; Rd = im12 - Rn$; Rn – op2 CMP Rn, <op2> sets the NZVC bits CMN Rn, <op2>; Rn - (-op2) sets the NZVC bits ; Rd = Rn * Rm MUL{S} {Rd,} Rn, Rm signed or unsigned Rd, Rn, Rm, Ra; Rd = Ra + Rn*RmMLA signed or unsigned MLS Rd, Rn, Rm, Ra ; Rd = Ra - Rn*Rmsigned or unsigned UDIV {Rd,} Rn, Rm ; Rd = Rn/Rmunsigned SDTV {Rd,} Rn, Rm ; Rd = Rn/Rmsigned Notes Ra Rd Rm Rn Rt represent 32-bit registers any 32-bit value: signed, unsigned, or address value if S is present, instruction will set condition codes {S} #im12 any value from 0 to 4095 #im16 any value from 0 to 65535 if Rd is present Rd is destination, otherwise Rn {Rd, } #n any value from 0 to 31 #off any value from -255 to 4095 any address within the ROM of the microcontroller label the value generated by <op2> op2 Examples of flexible operand **<op2>** creating the 32-bit number. E.g., **Rd = Rn+op2** ADD Rd, Rn, Rm ; op2 = RmADD Rd, Rn, Rm, LSL #n ; op2 = Rm<<n Rm is signed, unsigned ADD Rd, Rn, Rm, LSR #n ; op2 = Rm>>n Rm is unsigned ADD Rd, Rn, Rm, ASR #n ; op2 = Rm>>n Rm is signed ADD Rd, Rn, #constant ; op2 = constant, where X and Y are hexadecimal digits: • produced by shifting an 8-bit unsigned value left by any number of bits in the form **0x00XY00XY** • in the form **0xXY00XY00** in the form **0xXYXYXYX** R0 0x0000.0000 R1 256k Flash R2 **Condition code bits** ROM R3 0x0003.FFFF N negative R4 R5 General Z zero 0x2000.0000 purpose -R6 32k RAM V signed overflow **R**7 registers C carry or 0x2000.7FFF R8 R9 unsigned overflow 0x4000.0000 R10 I/O ports R11 R12 0x400F.FFFF Stack pointer R13 (MSP) Link register R14 (LR) 0xE000.0000 Program counter R15 (PC) Internal I/O PPB 0xE004.1FFF