

**Final Exam****Date:** May 13th 2016Printed Name: \_\_\_\_\_  
Last, First

Your signature is your promise that you have not cheated and will not cheat on this exam, nor will you help others to cheat on this exam. You will not reveal the contents of this exam to others who are taking the makeup thereby giving them an undue advantage:

Signature: \_\_\_\_\_

**Instructions:**

- Closed book and closed notes. No books, no papers, no data sheets (other than the last four pages of this Exam)
- No devices other than pencil, pen, eraser (no calculators, no electronic devices), please turn cell phones off.
- Please be sure that your answers to all questions (and all supporting work that is required) are contained in the space (boxes) provided. *Anything outside the boxes will be ignored in grading.*
- You have 180 minutes, so allocate your time accordingly.
- For all questions, unless otherwise stated, find the most efficient (time, resources) solution.
- Unless otherwise stated, make all I/O accesses friendly.
- *Please read the entire exam before starting. See supplement pages for Device I/O registers.*

<b>Problem 1</b>	10	
<b>Problem 2</b>	15	
<b>Problem 3</b>	10	
<b>Problem 4</b>	15	
<b>Problem 5</b>	15	
<b>Problem 6</b>	10	
<b>Problem 7</b>	15	
<b>Problem 8</b>	10	
<b>Total</b>	100	

(10) Question 1. Please place **one word or short phrase** that best answers each question in the box.

(1) Part a) Which registers are NOT pushed on the stack when an interrupt occurs?

R4-R11,R13/SP

(1) Part b) Why did we use the 7406 driver for interfacing the LED? .....

To increase current  
Needed > 8mA  
Protection

(1) Part c) If the average rate at which you call **Fifo\_Put** to enter data into a FIFO is less than the average rate at which you call **Fifo\_Get** to remove data, ..... could the FIFO ever become full?

Yes, there can be bursts of input

(1) Part d) Why do we add **static** to an otherwise local variable? ...

Persistence  
Allocated permanent RAM

(1) Part e) Why do we add **const** to an otherwise global variable? ...

Put in ROM, can't change it

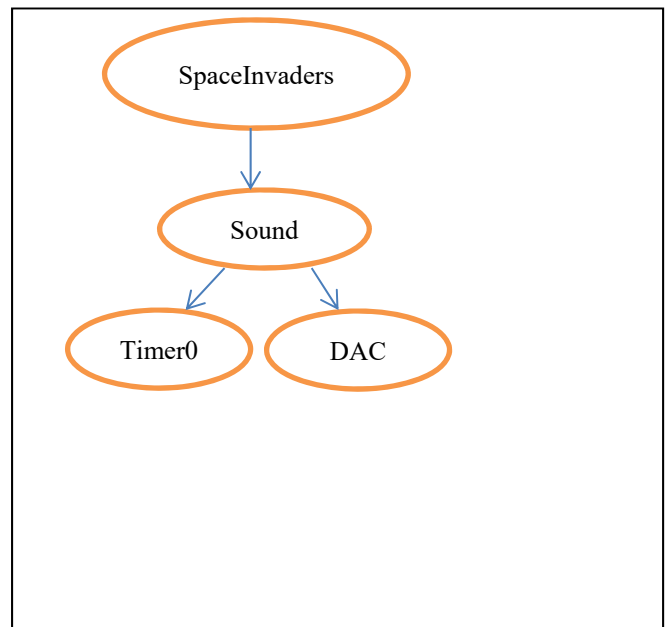
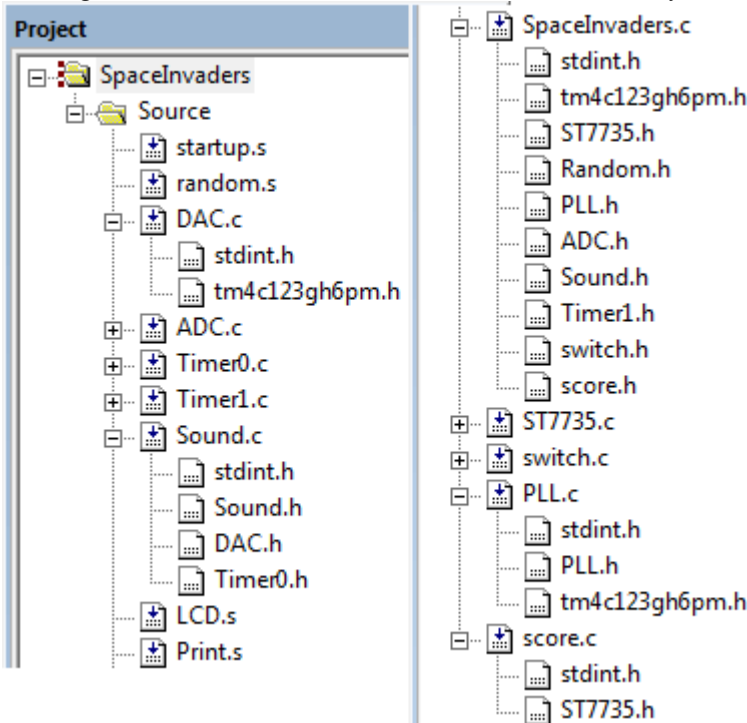
(1) Part f) Why do we add **static** to an otherwise global variable? ...

Reduce scope  
Private to file

(1) Part g) Why would you ever wish to use the PLL and slow down ... the bus clock so the software runs slower?

Save power  
Make battery last longer

(3) Part h) The following is the project window for my Lab 10 game. Draw a subset of the **call graph** of this system showing all the arrows into and out of the **Sound** module. Put your answer in the box.



**(15) Question 2: SysTick and Fixed-point**

**(6) Part a)** Write a function in C or assembly that uses SysTick delay for 1 millisecond. Assume that the SysTick is already initialized with `NVIC_ST_RELOAD=0x00FFFFFF`; `NVIC_ST_CTRL=0x05`; Assume that the clock is running at 50 MHz, such that each bus cycle is 20 ns.

```
void Delay1ms () {
    NVIC_ST_RELOAD = 50000; //10^-3 * 50x10^6 = 50x10^3
    NVIC_ST_CURRENT = 0;    // clear count
    // 1ms = 20ns*RELOAD, you could use any value from 49980 to 50000
    while (NVIC_ST_CTRL&0x00010000 == 0){}
}

void Delay1ms(void){ volatile uint32_t elapsedTime;
    uint32_t startTime = NVIC_ST_CURRENT_R;
    do{
        elapsedTime = (startTime-NVIC_ST_CURRENT_R)&0x00FFFFFF;
    }
    while(elapsedTime <= 50000); // use any value from 49980 to 50000
}
```

**(6) Part b)** Write C code initialization that initializes SysTick to interrupt every 100 us. In this initialization, in addition to enabling and arming SysTick interrupts you should also enable interrupts in the processor. Run at priority level 1. Assume the bus clock is 80 MHz. You do not need to write the SysTick ISR.

```
void SysTick_Init100us (void) {
    NVIC_ST_CTRL_R = 0;           // disable SysTick during setup
    NVIC_ST_RELOAD_R = 7999;      // 12.5*(7999+1)=100us
    NVIC_SYS_PRI3_R=(NVIC_SYS_PRI3_R&0x00FFFFFF)|0x20000000; // 1
    NVIC_ST_CURRENT_R = 0;       // clear flag
    NVIC_ST_CTRL_R = 0x07;       // arm, enable, source
    EnableInterrupts ();
}
```

**(3) Part c)** Consider the use of decimal fixed-point representation with a resolution of 0.01 cm. What is the area of a rectangle in  $\text{cm}^2$  whose width and length are represented by fixed-point integer values 250 and 110 respectively?

Width is represented as 250, Length as 110. The area is therefore  $2.50 * 1.10 = 2.75 \text{cm}^2$

**(10) Question 3: *FSM***

**(5) Part a)** A Moore FSM has 10 states, a 3-bit input (PB5-7), a 5-bit output (PB0-4) and a state dwell (or wait) time that can vary between 300ms to 1800ms. Complete the missing pieces in the definition of the struct for the State and FSM array declaration needed to implement the FSM.

```
#define Init 0
struct State{
    uint8_t out;        // ok if 16 or 32 bits
    unit16_t wait;     // ok if 32 bits
    unit8_t next[8];  // ok if 16 or 32 bits
}
```

```
typedef struct State State_t;
```

```
State_t FSM[ 10 ] = { ... contents are defined for you ...}
```

```
uint8_t CS; // Index into the FSM array indicating current state
```

**(5) Part a)** Complete the FSM engine loop inside the main below. Assume you are given a function `Delay1ms(uint32_t count)` that delays for count milliseconds.

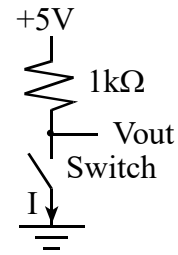
```
int main(void){
    PORTB_Init(); // Port B Initialization is done for you
    CS = Init;
    while(1){

        GPIO_PORTB_DATA_R = FSM[CS].out;
        Delay1ms(FSM[CS].wait);
        unit8_t in = (GPIO_PORTB_DATA_R & 0xE0) >> 5;
        CS = FSM[CS].next[in];

    }
}
```

**(15) Question 4: Interfacing**

**(3) Part a)** Consider this negative logic switch circuit used in a +5V digital system. Do not consider the switch to be ideal. Rather, assume the resistance of the switch when the switch is open is 1 MΩ, and the resistance of the switch when the switch is closed is 1 Ω. How much current flows in mA through the switch when the switch is not pressed?



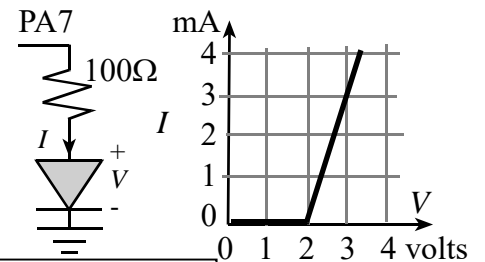
$$I = V/R = 5V/(1M+1k) = 5/1001 \text{ mA} \approx 0.005 \text{ mA}$$

How much current in mA flows through the switch when the switch is pressed?

$$I = V/R = 5V/(1+1k) = 5/1.001 \text{ mA} \approx 5 \text{ mA}$$

**(6) Part b)** The output on PA7 controls this LED. For LED voltages less than 2 volts, the LED current is 0. Assume the output high voltage of PA7 is 3.3V. For voltages above 2 volts, the LED current is

$$I = 3 * (V - 2), \text{ where } I \text{ is in mA, } V \text{ is in volts.}$$



What are the **current, voltage, and power** to the LED when it is on?

Let  $V$  be the LED voltage, and  $I$  be the current in both LED and resistor  
 The resistor voltage is  $(3.3 - V)$   
 The resistor current is  $(3.3 - V)/100$  in amps  
 The resistor current in mA is  $I = 33 - 10 * V$   
 The LED current equals the resistor current  
 $3 * V - 6 = 33 - 10 * V, 13 * V = 39, V = 3 \text{ volts, so } I = 3 \text{ mA}$   
 $P = V * I = 3 * 3 = 9 \text{ mW}$

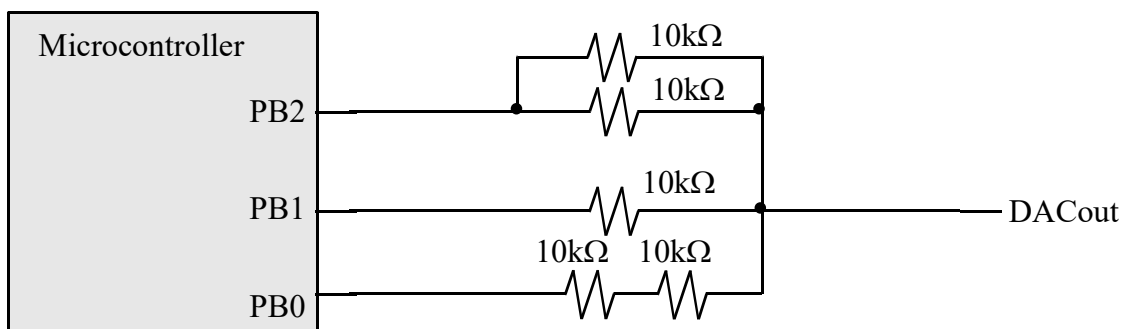
$V = 3 \text{ volts}$

$I = 3 \text{ mA}$

$P = 9 \text{ mW}$

**(6) Part c)** Design a 3-bit DAC using multiple 10k resistors. No values other than 10k are allowed.

Two 10k resistors in parallel are 5k, two 10k resistors in series are 20k. Ratio is 1/2/4, with the largest resistance at the least significant bit. PB0 has 20k, PB1 has 10k, PB2 has 5k.



**(15) Question 5: *FIFO queue***

You are asked to implement a FIFO that can handle up to **SIZE** elements. You cannot add additional the global variables. You cannot change the function prototypes.

**(3) Part a)** Write the routine that initializes the FIFO.

```
#define SIZE 10
uint8_t FIFO[SIZE], GetI, PutI, Count;
void Fifo_Init(void) { // Initialize FIFO

    Count = GetI = PutI = 0;
}
```

**(6) Part b)** Write the routine that puts data into the FIFO. If the FIFO is full, this routine should spin (i.e., wait) until there is room in FIFO for the data. You can add local variables.

```
// enter one byte into the FIFO
void Fifo_Put(uint8_t data) {
    while (Count == SIZE) {}
    Count++;
    FIFO[PutI] = data;
    PutI = (PutI+1)%SIZE;
}
```

**(6) Part c)** Write the routine that gets data from the FIFO. If the FIFO is empty, this routine should spin (i.e., wait) until there is data in FIFO to return. You can add local variables.

```
// return one byte of data
// if empty spin until there is data in FIFO
uint8_t Fifo_Get(void) {
    uint8_t data;
    while (Count == 0) {}
    data = FIFO[GetI];
    Count--;
    GetI = (GetI+1)%SIZE;
    return (data);
}
```

**(10) Question 6: *Local variables*****(3) Part a)** What does `LCD_OutDec` print to screen?

```

void add1(uint32_t *in){
    *in = (*in)+1;
}
void func(void){
    uint32_t a = 10;
    add1(&a);
    LCD_OutDec(a);
}

```

Answer (select one option):

- a) 11
- b) 10
- c) Some address of the stack region
- d) Address of `a`
- e) Unknown, unable to determine

a) 11

**(7) Part b)** Implement both subroutines `func` and `add1` in assembly. For the subroutine `func`, use binding, allocation, access and deallocation of the local variable `a` on the stack. Implement each line of C explicitly in assembly. Use AAPCS. You must clearly identify each of the different stages. Use C statements as comments.

```

add1
    LDR R1, [R0]
    ADD R1, #1 ; *in = (*in)+1
    STR R1, [R0]
    BX LR

a    equ 0
func
    PUSH {LR}
    SUB SP, #4 ; allocate a
    MOV R0, #10
    STR R0, [SP, #a] ; a=10
    MOV R0, SP
    BL add1 ; add1(&a);
    LDR R0, [SP, #a] ; a
    BL LCD_OutDec ; LCD_OutDec(a)
    ADD SP, #4 ; deallocate
    POP {PC}

```

**(15) Question 7: UART and ADC**

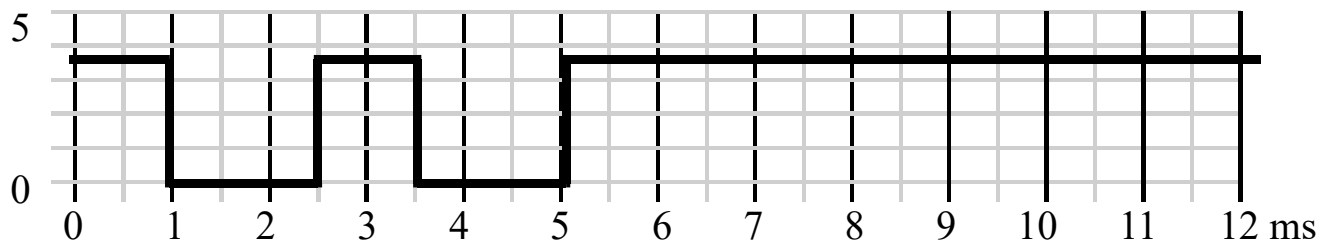
**(3) Part a)** A programmer set the `UART0_IBRD_R` to 50 and `UART0_FBRD_R` to 0. If the bus clock frequency were 80 MHz, what would be the baud rate? Pick one answer.

- i. 80 kbps
- ii. 100 kbps
- iii. 1 Mbps
- iv. 120 kbps
- v. 16 kbps
- vi. None of the above

$$\text{ii. } 80,000,000/16/50 = 100,000 \text{ bps}$$

**(6) Part b)** Assume a serial port operating with a baud rate of 2000 bits per second. Draw the UART waveform when the decimal value 140 is transmitted. You may assume the channel is idle before and after the frame. Time flows from left to right. The dark vertical black lines correspond to 1-ms boundaries. Assume the frame begins at time = 1 ms, and show the waveform from 0 to 12 ms.

$$140 = 128 + 8 + 4, \text{ binary} = 10001100$$



**(3) Part c)** You have a 12-bit, 0 to 3V range ADC. If the ADC input is 1.25 V, what will be the digital value **in hex** returned by this ADC?

$$4095 * 1.25/3 = 1706 = 0x6AA$$

0x6AA

**(3) Part d)** You are attempting to capture a sinusoidal sound with a frequency of 8 kHz. The `ADC0_PC_R` is set to 0001, which supports a maximum of 125k samples/sec. Using the 12-bit ADC and periodic interrupt, you have programmed the SysTick to interrupt at a frequency of 12 kHz. During the SysTick ISR you collect one ADC sample. Is it possible to recreate the original signal from the captured samples? If your answer is *yes*, explain how. If your answer is *no*, what is the term used to refer to this loss of information?

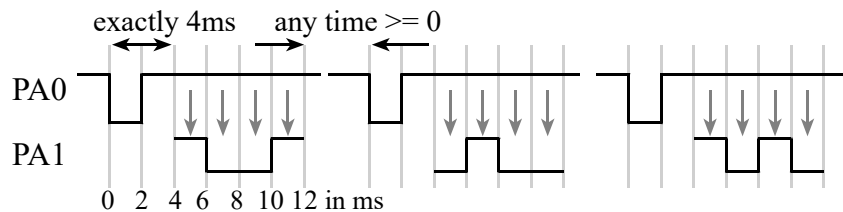
No, The sampling rate in this case that satisfies the Nyquist theorem is 16kHz (2\*8kHz). The chosen sampling frequency of 12 kHz is insufficient and will cause **aliasing**.



**(10) Question 8: Design Problem**

You are to implement one-directional communication between two microcontrollers (sender and receiver) using two pins PA0 and PA1. The sender microcontroller programs these pins as output and the receiver programs them as input. PA0 is used as the control and PA1 is used to transfer the actual data. The bit-level protocol is as follows: when it has data to transmit, the sender sends a pulse on PA0, which is a low for 2 ms, followed by a high for 2 ms. The first bit of transmission (on PA1) immediately follows, with each bit sent lasting for exactly 2 ms. Each 4-bit transmission is preceded by a pulse on PA0. The time between transmissions can be any value greater than or equal to zero. See in the timeline below that the first byte transferred is 0x29. Note that each byte of data is transmitted as two 4-bit nibbles with the bit order as 0,1,2,3 then 4,5,6,7. The vertical arrows mark when you should read the input data. You may assume PA1 and PA0 are initialized to inputs.

**(8) Part a)** The overall goal of the communication is to transfer data from the sender to the receiver. In this section you will write a function that receives one byte of data. You may assume that you are given a function `Delay1ms()`.



You may assume the software execution time is negligible compared to the Delay function

```
uint8_t ReceiveChar(void) {
    uint8_t data;
    while((GPIO_PORTA_DATA_R&0x01)==1) {}
    while((GPIO_PORTA_DATA_R&0x01)==0) {}
    Delay1ms(); Delay1ms(); Delay1ms();
    data = (GPIO_PORTA_DATA_R&0x02)>>1; Delay1ms(); Delay1ms();
    data = data|(GPIO_PORTA_DATA_R&0x02); Delay1ms(); Delay1ms();
    data = data|((GPIO_PORTA_DATA_R&0x02)<<1); Delay1ms(); Delay1ms();
    data = data|((GPIO_PORTA_DATA_R&0x02)<<2);
    while((GPIO_PORTA_DATA_R&0x01)==1) {}
    while((GPIO_PORTA_DATA_R&0x01)==0) {}
    Delay1ms(); Delay1ms(); Delay1ms();
    data = data|((GPIO_PORTA_DATA_R&0x02)<<3); Delay1ms(); Delay1ms();
    data = data|((GPIO_PORTA_DATA_R&0x02)<<4); Delay1ms(); Delay1ms();
    data = data|((GPIO_PORTA_DATA_R&0x02)<<5); Delay1ms(); Delay1ms();
    data = data|((GPIO_PORTA_DATA_R&0x02)<<6);
    return data;
}
```

**(2) Part b)** What is maximum achievable **bandwidth** in bits/sec for this communication scenario?

$4\text{bits}/12\text{ms} = 4000\text{bits}/12\text{s} = 1000/3 = 333\text{bits/sec}$ ,  $4/6 * (500\text{bps}) = 333\text{bps}$

**Memory access instructions**

```

LDR   Rd, [Rn]           ; load 32-bit number at [Rn] to Rd
LDR   Rd, [Rn,#off]     ; load 32-bit number at [Rn+off] to Rd
LDR   Rd, =value        ; set Rd equal to any 32-bit value (PC rel)
LDRH  Rd, [Rn]           ; load unsigned 16-bit at [Rn] to Rd
LDRH  Rd, [Rn,#off]     ; load unsigned 16-bit at [Rn+off] to Rd
LDRSH Rd, [Rn]           ; load signed 16-bit at [Rn] to Rd
LDRSH Rd, [Rn,#off]     ; load signed 16-bit at [Rn+off] to Rd
LDRB  Rd, [Rn]           ; load unsigned 8-bit at [Rn] to Rd
LDRB  Rd, [Rn,#off]     ; load unsigned 8-bit at [Rn+off] to Rd
LDRSB Rd, [Rn]           ; load signed 8-bit at [Rn] to Rd
LDRSB Rd, [Rn,#off]     ; load signed 8-bit at [Rn+off] to Rd
STR   Rt, [Rn]           ; store 32-bit Rt to [Rn]
STR   Rt, [Rn,#off]     ; store 32-bit Rt to [Rn+off]
STRH  Rt, [Rn]           ; store least sig. 16-bit Rt to [Rn]
STRH  Rt, [Rn,#off]     ; store least sig. 16-bit Rt to [Rn+off]
STRB  Rt, [Rn]           ; store least sig. 8-bit Rt to [Rn]
STRB  Rt, [Rn,#off]     ; store least sig. 8-bit Rt to [Rn+off]
PUSH  {Rt}               ; push 32-bit Rt onto stack
POP   {Rd}               ; pop 32-bit number from stack into Rd
ADR   Rd, label          ; set Rd equal to the address at label
MOV{S} Rd, <op2>         ; set Rd equal to op2
MOV   Rd, #im16          ; set Rd equal to im16, im16 is 0 to 65535
MVN{S} Rd, <op2>        ; set Rd equal to -op2

```

**Branch instructions**

```

B     label ; branch to label      Always
BEQ  label ; branch if Z == 1     Equal
BNE  label ; branch if Z == 0     Not equal
BCS  label ; branch if C == 1     Higher or same, unsigned ≥
BHS  label ; branch if C == 1     Higher or same, unsigned ≥
BCC  label ; branch if C == 0     Lower, unsigned <
BLO  label ; branch if C == 0     Lower, unsigned <
BMI  label ; branch if N == 1     Negative
BPL  label ; branch if N == 0     Positive or zero
BVS  label ; branch if V == 1     Overflow
BVC  label ; branch if V == 0     No overflow
BHI  label ; branch if C==1 and Z==0 Higher, unsigned >
BLS  label ; branch if C==0 or Z==1 Lower or same, unsigned ≤
BGE  label ; branch if N == V     Greater than or equal, signed ≥
BLT  label ; branch if N != V     Less than, signed <
BGT  label ; branch if Z==0 and N==V Greater than, signed >
BLE  label ; branch if Z==1 or N!=V Less than or equal, signed ≤
BX   Rm      ; branch indirect to location specified by Rm
BL   label   ; branch to subroutine at label
BLX  Rm      ; branch to subroutine indirect specified by Rm

```

**Interrupt instructions**

```

CPSIE I           ; enable interrupts (I=0)
CPSID I           ; disable interrupts (I=1)

```

**Logical instructions**

```

AND{S} {Rd,} Rn, <op2> ; Rd=Rn&op2      (op2 is 32 bits)
ORR{S} {Rd,} Rn, <op2> ; Rd=Rn|op2      (op2 is 32 bits)
EOR{S} {Rd,} Rn, <op2> ; Rd=Rn^op2      (op2 is 32 bits)
BIC{S} {Rd,} Rn, <op2> ; Rd=Rn&(~op2) (op2 is 32 bits)
ORN{S} {Rd,} Rn, <op2> ; Rd=Rn|(~op2) (op2 is 32 bits)
LSR{S} Rd, Rm, Rs      ; logical shift right Rd=Rm>>Rs (unsigned)
LSR{S} Rd, Rm, #n      ; logical shift right Rd=Rm>>n (unsigned)
ASR{S} Rd, Rm, Rs      ; arithmetic shift right Rd=Rm>>Rs (signed)

```

```
ASR{S} Rd, Rm, #n ; arithmetic shift right Rd=Rm>>n (signed)
LSL{S} Rd, Rm, Rs ; shift left Rd=Rm<<Rs (signed, unsigned)
LSL{S} Rd, Rm, #n ; shift left Rd=Rm<<n (signed, unsigned)
```

**Arithmetic instructions**

```
ADD{S} {Rd,} Rn, <op2> ; Rd = Rn + op2
ADD{S} {Rd,} Rn, #im12 ; Rd = Rn + im12, im12 is 0 to 4095
SUB{S} {Rd,} Rn, <op2> ; Rd = Rn - op2
SUB{S} {Rd,} Rn, #im12 ; Rd = Rn - im12, im12 is 0 to 4095
RSB{S} {Rd,} Rn, <op2> ; Rd = op2 - Rn
RSB{S} {Rd,} Rn, #im12 ; Rd = im12 - Rn
CMP Rn, <op2> ; Rn - op2 sets the NZVC bits
CMN Rn, <op2> ; Rn - (-op2) sets the NZVC bits
MUL{S} {Rd,} Rn, Rm ; Rd = Rn * Rm signed or unsigned
MLA Rd, Rn, Rm, Ra ; Rd = Ra + Rn*Rm signed or unsigned
MLS Rd, Rn, Rm, Ra ; Rd = Ra - Rn*Rm signed or unsigned
UDIV {Rd,} Rn, Rm ; Rd = Rn/Rm unsigned
SDIV {Rd,} Rn, Rm ; Rd = Rn/Rm signed
```

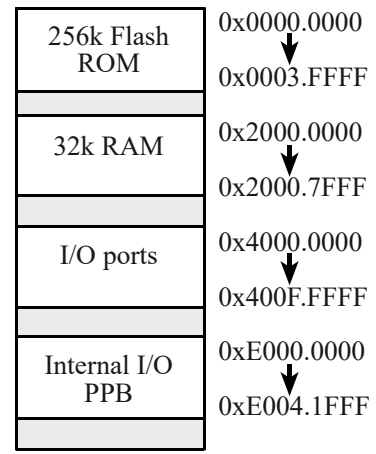
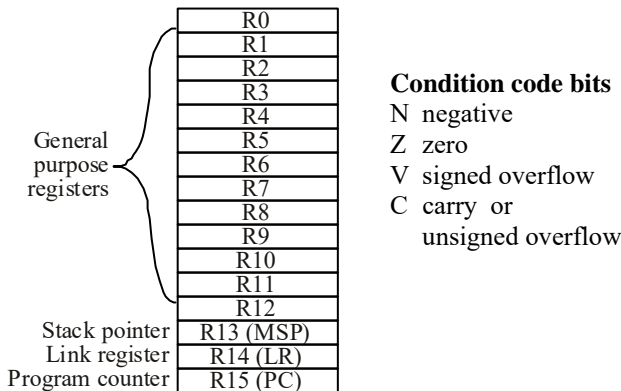
**Notes** Ra Rd Rm Rn Rt represent 32-bit registers

```
value any 32-bit value: signed, unsigned, or address
{S} if S is present, instruction will set condition codes
#im12 any value from 0 to 4095
#im16 any value from 0 to 65535
{Rd,} if Rd is present Rd is destination, otherwise Rn
#n any value from 0 to 31
#off any value from -255 to 4095
label any address within the ROM of the microcontroller
op2 the value generated by <op2>
```

**Examples of flexible operand <op2> creating the 32-bit number. E.g., Rd = Rn+op2**

```
ADD Rd, Rn, Rm ; op2 = Rm
ADD Rd, Rn, Rm, LSL #n ; op2 = Rm<<n Rm is signed, unsigned
ADD Rd, Rn, Rm, LSR #n ; op2 = Rm>>n Rm is unsigned
ADD Rd, Rn, Rm, ASR #n ; op2 = Rm>>n Rm is signed
ADD Rd, Rn, #constant ; op2 = constant, where X and Y are hexadecimal digits:
```

- produced by shifting an 8-bit unsigned value left by any number of bits
- in the form 0x00XY00XY
- in the form 0xXY00XY00
- in the form 0xXYXYXYXY



```
DCB 1,2,3 ; allocates three 8-bit byte(s)
DCW 1,2,3 ; allocates three 16-bit halfwords
DCD 1,2,3 ; allocates three 32-bit words
SPACE 4 ; reserves 4 bytes
```

Address	7	6	5	4	3	2	1	0	Name
\$400F.E608			GPIOF	GPIOE	GPIOD	GPIOC	GPIOB	GPIOA	SYSCTL RCGCGPIO R
\$4000.53FC	DATA	DATA	DATA	DATA	DATA	DATA	DATA	DATA	GPIO PORTB DATA R
\$4000.5400	DIR	DIR	DIR	DIR	DIR	DIR	DIR	DIR	GPIO PORTB DIR R
\$4000.5420	SEL	SEL	SEL	SEL	SEL	SEL	SEL	SEL	GPIO PORTB AFSEL R
\$4000.551C	DEN	DEN	DEN	DEN	DEN	DEN	DEN	DEN	GPIO PORTB DEN R

Table 4.5. TM4C123 Port B parallel ports. Each register is 32 bits wide. Bits 31 – 8 are zero.

Address	31	30	29-7	6	5	4	3	2	1	0	Name
0xE000E100		F	...	UART1	UART0	E	D	C	B	A	NVIC EN0 R

Address	31-24	23-17	16	15-3	2	1	0	Name
\$E000E010	0	0	COUNT	0	CLK_SRC	INTEN	ENABLE	NVIC ST CTRL R
\$E000E014	0	24-bit RELOAD value						NVIC ST RELOAD R
\$E000E018	0	24-bit CURRENT value of SysTick counter						NVIC ST CURRENT R

Address	31-29	28-24	23-21	20-8	7-5	4-0	Name
\$E000ED20	SYSTICK	0	PENDSV	0	DEBUG	0	NVIC SYS PRI3 R

Table 9.6. SysTick registers.

Table 9.6 shows the SysTick registers used to create a periodic interrupt. SysTick has a 24-bit counter that decrements at the bus clock frequency. Let  $f_{BUS}$  be the frequency of the bus clock, and let  $n$  be the value of the **RELOAD** register. The frequency of the periodic interrupt will be  $f_{BUS}/(n+1)$ . First, we clear the **ENABLE** bit to turn off SysTick during initialization. Second, we set the **RELOAD** register. Third, we write to the **NVIC\_ST\_CURRENT\_R** value to clear the counter. Lastly, we write the desired mode to the control register, **NVIC\_ST\_CTRL\_R**. To turn on the SysTick, we set the **ENABLE** bit. We must set **CLK\_SRC**=1, because **CLK\_SRC**=0 external clock mode is not implemented. We set **INTEN** to arm SysTick interrupts. The standard name for the SysTick ISR is **SysTick\_Handler**.

Address	31-2			1			0			Name	
\$400F.E638				ADC1			ADC0			SYSCTL_RCGCADC R	
	31-14	13-12	11-10	9-8	7-6	5-4	3-2	1-0			
\$4003.8020		SS3		SS2		SS1		SS0	ADC0_SSPRI R		
	31-16			15-12		11-8		7-4		3-0	
\$4003.8014				EM3		EM2		EM1		EM0	ADC0_EMUX R
	31-4			3		2		1		0	
\$4003.8000				ASEN3		ASEN2		ASEN1		ASEN0	ADC0_ACTSS R
\$4003.80A0				MUX0						ADC0_SSMUX3 R	
\$4003.80A4				TS0		IE0		END0		D0	ADC0_SSCTL3 R
\$4003.8028				SS3		SS2		SS1		SS0	ADC0_PSSI R
\$4003.8004				INR3		INR2		INR1		INR0	ADC0_RIS R
\$4003.8008				MASK3		MASK2		MASK1		MASK0	ADC0_IM R
\$4003.8FC4				Speed						ADC0_PC R	
	31-12			11-0							
\$4003.80A8				DATA						ADC0_SSFIFO3 R	

Table 10.3. The TM4C ADC registers. Each register is 32 bits wide. LM3S has 10-bit data.

Set Speed to 0001 for slow speed operation. The ADC has four sequencers, but we will use only sequencer 3. We set the **ADC\_SSPRI\_R** register to 0x3210 to make sequencer 3 the lowest priority. Because we are using just one sequencer, we just need to make sure each sequencer has a unique priority. We set bits 15–12 (**EM3**) in the **ADC\_EMUX\_R** register to specify how the ADC will be triggered. If we specify software start (**EM3**=0x0), then the software writes an 8 (**SS3**) to the **ADC\_PSSI\_R** to initiate a conversion on sequencer 3. Bit 3 (**INR3**) in the **ADC\_RIS\_R** register will be set when the conversion is complete. We can enable and disable the sequencers using the **ADC\_ACTSS\_R** register. Which channel we sample is configured by writing to the **ADC\_SSMUX3\_R** register. The **ADC\_SSCTL3\_R** register specifies the mode of the ADC sample. Clear **TS0**. We set **IE0** so that the **INR3** bit is set on ADC conversion, and clear it when no flags are needed. We will set **IE0** for both interrupt and busy-wait synchronization. When using sequencer 3, there is only one sample, so **END0** will always be set, signifying this sample is the end of the sequence. Clear the **D0** bit. The **ADC\_RIS\_R**

register has flags that are set when the conversion is complete, assuming the **IE0** bit is set. Do not set bits in the **ADC\_IM\_R** register because we do not want interrupts. Write one to **ADC\_ISC\_R** to clear the corresponding bit in the **ADC\_RIS\_R** register.

UART0 pins are on PA1 (transmit) and PA0 (receive). The **UART0\_IBRD\_R** and **UART0\_FBRD\_R** registers specify the baud rate. The baud rate **divider** is a 22-bit binary fixed-point value with a resolution of  $2^{-6}$ . The **Baud16** clock is created from the system bus clock, with a frequency of (Bus clock frequency)/**divider**. The baud rate is

$$\text{Baud rate} = \text{Baud16}/16 = (\text{Bus clock frequency})/(16 * \text{divider})$$

We set bit 4 of the **UART0\_LCRH\_R** to enable the hardware FIFOs. We set both bits 5 and 6 of the **UART0\_LCRH\_R** to establish an 8-bit data frame. The **RTRIS** is set on a receiver timeout, which is when the receiver FIFO is not empty and no incoming frames have occurred in a 32-bit time period. The arm bits are in the **UART0\_IM\_R** register. To acknowledge an interrupt (make the trigger flag become zero), software writes a 1 to the corresponding bit in the **UART0\_IC\_R** register.

We set bit 0 of the **UART0\_CTL\_R** to enable the UART. Writing to **UART0\_DR\_R** register will output on the UART. This data is placed in a 16-deep transmit hardware FIFO. Data are transmitted first come first serve. Received data are place in a 16-deep receive hardware FIFO. Reading from **UART0\_DR\_R** register will get one data from the receive hardware FIFO. The status of the two FIFOs can be seen in the **UART0\_FR\_R** register (FF is FIFO full, FE is FIFO empty). The standard name for the UART0 ISR is **UART0\_Handler**. **RXIFLSEL** specifies the receive FIFO level that causes an interrupt (010 means interrupt on  $\geq \frac{1}{2}$  full, or 7 to 8 characters). **TXIFLSEL** specifies the transmit FIFO level that causes an interrupt (010 means interrupt on  $\leq \frac{1}{2}$  full, or 9 to 8 characters).

\$4000.C000	31-12	11	10	9	8	7-0		Name	
		OE	BE	PE	FE	DATA		UART0_DR_R	
\$4000.C004	31-3			3	2	1	0	UART0_RSR_R	
		OE	BE	PE	FE				
\$4000.C018	31-8	7	6	5	4	3	2-0	UART0_FR_R	
		TXFE	RXFF	TXFF	RXFE	BUSY			
\$4000.C024	31-16	15-0						UART0_IBRD_R	
		DIVINT							
\$4000.C028	31-6				5-0				UART0_FBRD_R
					DIVFRAC				
\$4000.C02C	31-8	7	6-5	4	3	2	1	0	UART0_LCRH_R
		SPS	WPEN	FEN	STP2	EPS	PEN	BRK	
\$4000.C030	31-10	9	8	7	6-3	2	1	0	UART0_CTL_R
		RXE	TXE	LBE		SIRLP	SIREN	UARTEN	
\$4000.C034	31-6			5-3			2-0		UART0_IFLS_R
				RXIFLSEL			TXIFLSEL		
\$4000.C038	31-11	10	9	8	7	6	5	4	UART0_IM_R
		OEIM	BEIM	PEIM	FEIM	RTIM	TXIM	RXIM	
\$4000.C03C		OERIS	BERIS	PERIS	FERIS	RTRIS	TXRIS	RXRIS	UART0_RIS_R
\$4000.C040		OEMIS	BEMIS	PEMIS	FEMIS	RTMIS	TXMIS	RXMIS	UART0_MIS_R
\$4000.C044		OEIC	BEIC	PEIC	FEIC	RTIC	TXIC	RXIC	UART0_IC_R

Table 11.2. UART0 registers. Each register is 32 bits wide. Shaded bits are zero.