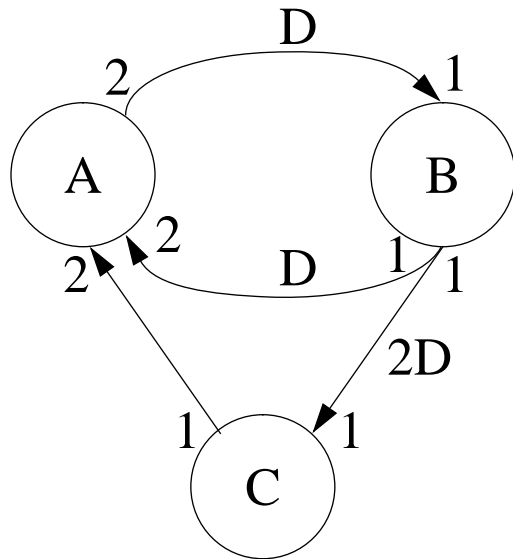


## Existence of Single Appearance Schedules

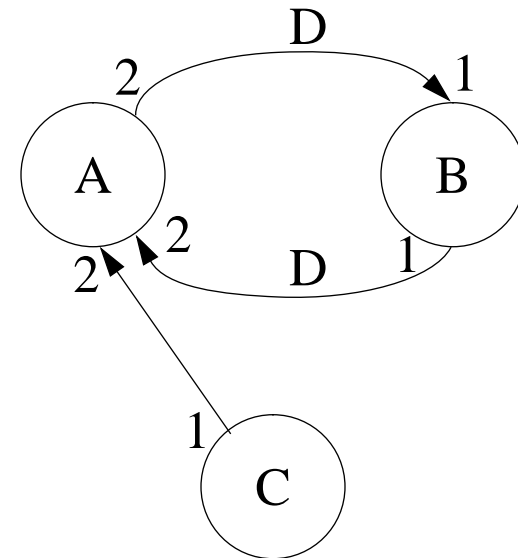
- Every acyclic SDF graph has a single appearance schedule.
- An arbitrary SDF graph has a single appearance schedule iff each strongly connected component has a single appearance schedule.
- A strongly connected SDF graph has a single appearance schedule only if the nodes can be partitioned into  $Z_1$  and  $Z_2$  such that for any arc  $\alpha$  directed from  $Z_1$  to  $Z_2$ ,  $delay(\alpha) \mathbf{q}_{\text{sink}(\alpha)} \times consumed(\alpha)$ .
- “Subindependent Partition”

## Finding a Subindependent Partition

Remove each arc  $\alpha$  for which  
 $delay(\alpha) \geq \mathbf{q}_{\text{sink}(\alpha)}$



Any “source” strongly connected component is subindependent of the rest of the graph.



{C} is subindependent of {A, B}

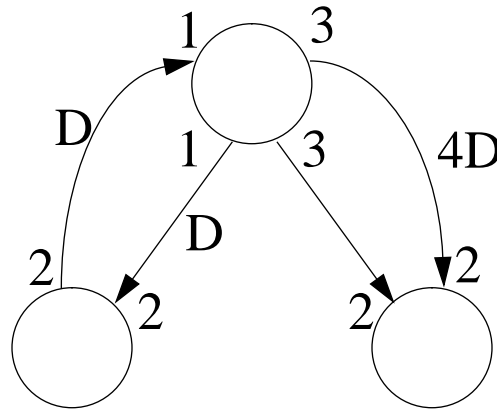
## Components of the Scheduling Framework

- Subindependent Partitioning Algorithm.
- Acyclic Scheduling Algorithm.
- Tight Scheduling Algorithm.

algorithm *schedule-loops*

- Consolidate each strongly connected component  $\{C_i\} \rightarrow$  *acyclic graph*  $G'$ .
- Apply acyclic scheduling algorithm to  $G' \rightarrow$  single appearance schedule of strongly connected components  $S'$ .
- For each strongly connected component  $C_i$ 
  - Apply the subindependence partitioning algorithm
  - If a subindependent partition  $(P_1, P_2)$  is found
    - Determine looping factors:  $r_i = gcd(\{\mathbf{q}_N \mid N \in P_i\})$ , for  $i = 1, 2$
    - Apply `schedule-loops()` to  $P_1$  and  $P_2 \rightarrow$  subschedules  $S_1$  and  $S_2$
    - Replace appearance of  $C_i$  in  $S'$  with  $(r_1 S_1) (r_2 S_2)$
  - If no subindependent partition is found
    - Apply the tight scheduling algorithm  $\rightarrow$  subschedule  $S_T$
    - Replace appearance of  $C_i$  in  $S'$  with  $S_T$

- We say that a strongly connected SDF graph is *tightly interdependent* if it does not have a subindependent partition.
- If  $Z_1$  and  $Z_2$  are subsets of nodes in an SDF graph such that  $Z_1 \cap Z_2 \neq \emptyset$  and  $subgraph(Z_1)$  &  $subgraph(Z_2)$  are both tightly interdependent, then  $subgraph(Z_1 \cup Z_2)$  is tightly interdependent.
- *Tightly interdependent components.*



## Properties of the Scheduling Framework

- Constructs a single appearance schedules whenever one exists.
- Actors outside the tightly interdependent components are scheduled with only one appearance.
- For actors inside tightly interdependent components, the number of appearances is determined entirely by the tight scheduling algorithm.