Department of Electrical and Computer Engineering The University of Texas at Austin

Name:____

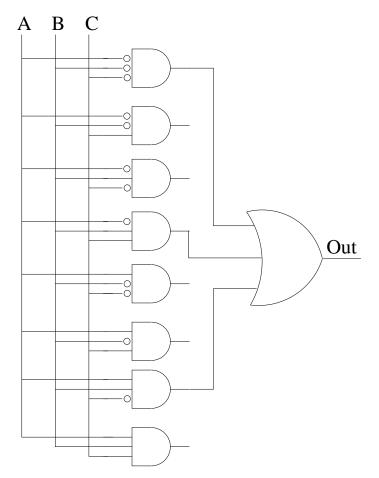
EE 306, Fall 2011 Yale Patt, Instructor Faruk Guvenilir, Milad Hashemi, Jennifer Davis, Garrett Galow, Ben Lin, Taylor Morrow, Stephen Pruett, Jee Ho Ryoo TAs Final Exam, December 9, 2011

	Part A:		
	Problem 1 (10 points):		
	Problem 2 (10 points):		
	Problem 3 (10 points):		
	Problem 4 (10 points):		
	Problem 5 (10 points):	Part A (50 points):	
	Part B:		
	Problem 6 (20 points):		
	Problem 7 (20 points):		
	Problem 8 (20 points):		
	Problem 9 (20 points):	Total (130 points):	
Note: Please be sure that your answer	ers to all questions (and all suppor	ting work that is required) are	contained in the space
Note: Please be sure your name is re	ecorded on each sheet of the exam	1.	
I will not cheat on this exam.			
Signature			

GOOD LUCK! (HAVE A GREAT SEMESTER BREAK) Name:

Problem 1. (10 points):

Part a. (5 points): Construct the output of the truth table for the PLA shown.



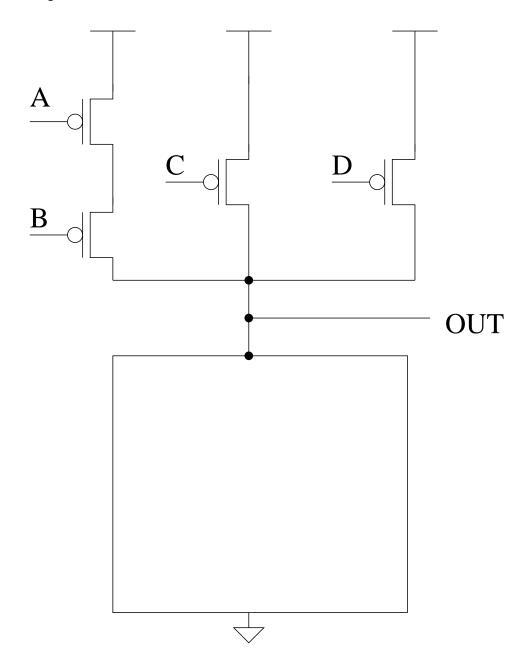
A	В	C	Out
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

Name:

Part b. (5 points): In the transistor circuit below, all transistors in the path to the power supply are shown. None of the transistors in the path to ground are shown.

Your job:

1. Draw the missing transistor circuit in the box.



Problem 2. (10 points): The following program is assembled and stored in the LC-3's memory. The PC is initially set to x3000. The program is run until the computer halts.

Your job: What is contained in location B after the computer stops?

	.ORIG	x3000
	AND	R0,R0,#0
	NOT	R1,R0
	ADD	R5,R0,#3
	ADD	R0,R0,#1
	ADD	R0,R0,R0
	ADD	R0,R0,R0
	ADD	R0,R0,R0
	NOT	R3,R0
	AND	R1,R3,R1
A	ADD	R0,R0,R0
	NOT	R3,R0
	AND	R1,R3,R1
	ADD	R5,R5,#-1
	BRp	A
	ST	R1,B
	TRAP	x25
В	.BLKW	1

What is the value in location B?

Name:			

Problem 3. (10 points): This problem involves a new 16-bit floating point data type, specified as follows:

S	ig	1	Exponent						Fr	act	io	n			
				l	l		ı	ı	l		ı	l	l	l	ı

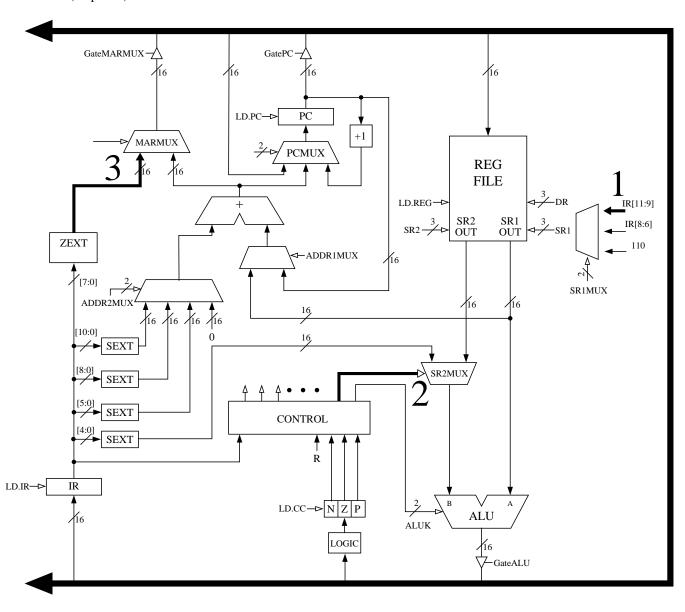
To add two floating point values, we first make sure their binary points line up (they have the same exponents).

The assembly program shown below, after the missing instructions have been filled in, compares the exponents of two floating point numbers that have been previously loaded into locations A and B. If the exponents are the same, R5 is set to 0 before the RET is taken. If the exponents are different, R5 is set to 1 before the RET is taken.

Your job: Fill in the missing instructions.

	.ORIG x3000 ST R0,SaveR0 ST R1,SaveR0
	LD R2, MASK AND R5, R5, #0 LD R0,A LD R1,B
	NOT R1, R1 ADD R1,R1,#1
DONE	BRZ DONE ADD R5,R5,#1 LD R0,SaveR0 LD R1,SaveR1
	RET
MASK	
А	.BLKW #1
B SaveR0	.BLKW #1 .BLKW #1
	.BLKW #1
С	
	.END

Problem 4. (10 points):



- 1. What opcodes use IR[11:9] as inputs to SR1?
- 2. Where does the control signal of this mux come from? Be specific!

3. What opcodes use this input to the MARMUX?

	.ORIG x3000	;	Line I
	LD R3, L2	;	2
	LDR R0, R3, #0	;	3
	LDR R1, R3, #1	;	4
	NOT R2, R1	;	5
	ADD R2, R2, #1	;	6
L1	ADD R0, R0, R2	;	7
	BRzp L1	;	8
	ADD R0, R0, R1	;	9
	ST RO, L3	;	10
	HALT	;	11
L2	.FILL x3100	;	12
L3	.FILL x3200	;	13
	.END	;	14
Part A Afte	or the instruction at line 6 has	ρv	ecuted, what are the contents of R0,R1,and R2? NOTE: the correct answer in
	one of the following: A, -A,		
cach case is	one of the following. A, -A,	υ,	D, 0, 1, -1.
R0:	R1:		R2:
_			
Part B. Then	re is a bug in the program. The	in	struction at line should be

Problem 5. (10 points): The modulo operator (A mod B) is the remainder one gets when dividing A by B. For example,

The program below is supposed to perform A mod B, where A is in x3100 and B is in x3101. The result should be stored at location x3200. However, the programmer made a serious mistake, so the program does not work. You can assume that

Name:_

10 mod 5 is 0, 12 mod 7 is 5.

A and B are both positive integers.



Problem 6. (20 points): A free list is a collection of blocks of consecutive memory locations of various sizes that are not being used by currently executing programs. A free list is normally organized as a linked list, where each element in the linked list is associated with a single block of memory. Each element consists of three words: the address of the next element in the linked list, the number of consecutive memory locations in this block, and the starting address of the block. R1 contains the address of a memory location that points to the first node in the free list.

R1:	xC000	xC000:	x8000	x8000:	xA000	xA000:	x0000
				x8001:	x0100	xA001:	x0010
				x8002:	x6000	xA002:	x7050

The free list above consists of two nodes, one of size x100 comprising M[x6000] to M[x60FF] and one of size x10 comprising locations M[x7050] to M[x705F].

A procedure MALLOC is used to provide blocks of storage to programs that request them.

If Program A needs n words of memory, it loads n into R2 and does a JSR to MALLOC. MALLOC finds the first block in the free list that can satisfy the request, loads the starting address of the block into R0, updates the free list to reflect the fact that those n words are no longer available, and does a JMP R7. If MALLOC can't find a block that can satisfy the request, x0000 is returned in R0. If the block that supplied the n-words consisted of exactly n-words (a perfect fit), then no words from that block are still available and so the node is removed from the free list.

On the next page is the procedure MALLOC. Your job: Add the missing instructions.

Name:	
MALLOC	ST R1, SAVE_R1
	ST R3, SAVE_R3
	ST R4, SAVE_R4
	ST R5, SAVE_R5
	AND R0, R0, #0
	NOT R3, R2
	ADD R3, R3, #1
NEXT_NODE	LDR R4, R1, #0
	BRz RETURN
	LDR R5, R4, #1
	ADD R5, R3, R5
	BRz PERFECT_FIT
	BRp FRAGMENT
-	
	BRnzp NEXT_NODE
PERFECT_FIT	LDR R0, R4, #2
	STR R4, R1, #0
	BRnzp RETURN
FRAGMENT	LDR R0, R4, #2
	STR R5, R4, #1
	STR R1, R4, #2
RETURN	LD R5, SAVE_R5
	LD R4, SAVE_R4
	LD R3, SAVE_R3
	LD R1, SAVE_R1
	RET
SAVE_R1	.BLKW 1
SAVE_R3	.BLKW 1
SAVE_R4	.BLKW 1
SAVE_R5	.BLKW 1

	er stops due to a breal	uring the processing of an LC-3 program by the data path we have been using in class, the kpoint set at x3000. The contents of certain registers and memory locations at that time are
R:	2 through R7: 2 M[x3000]: 2 M[x3003]: 2	x1263
In 15 of	those clock cycles, th	ecutes exactly four instructions. To accomplish this, a number of clock cycles are required. he bus must be utilized. The table below lists those 15 clock cycles in sequential order, along lonto the LC-3 bus in each.
	BUS	
1st:	x3000	
2nd:	x1263	
3rd:	x009A	
4th:	x3001	
5th:	xA000	
6th:		
7th:		
8th:		
9th:		
10th:		
11th:		
12th:		
13th:	x3003	
14th:	x1263	
15th:	x009D	
Part a:	Fill in the missing en	tries above.
Part b:	What are the four ins	structions that were executed?
Part c:	What are the contents	s of R0 and R1 after the four instructions execute?
	R0:	R1: 10

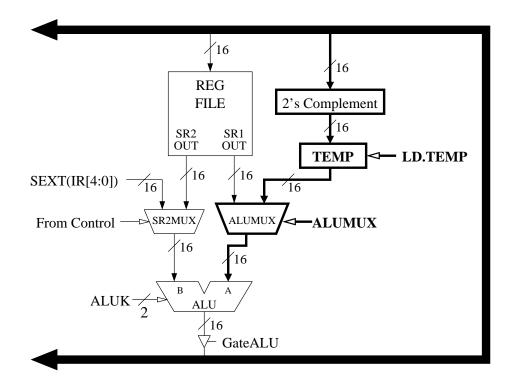
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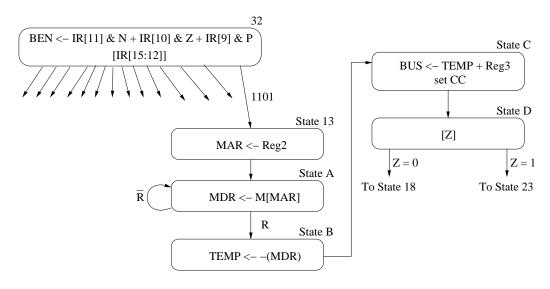
Problem 8. (20 points): Let's use the unused opcode to implement a new instruction, as shown below:

1101	1	Re	gĺ	Re	g2	00	0	Reg	3
15	12	11	9	8	6	5	3	2	0

To accomplish this, we will need a small addition to the data path, shown below in boldface:



The following five additional states are needed to control the data path to carry out the work of this instruction.



Note: State B loads the negative of the contents of MDR into TEMP.

Part a: Complete the table below by identifying the values of the control signals needed to carry out the work of each state.

Note: For a particular state, if the value of a control signal does not matter, fill it with an X.

	LD.PC	LD.MAR	LD.MDR	LD.CC	LD.TEMP	GatePC	GateMDR	GateALU	SP1MITY[1:0]	ALUMUX	10.1171111	ALUM[1:0]	MIO.EN	R.W
State 13														
State A														
State B														
State C														
State D														

LD.PC	0: load not enabled 1: load enabled	SR1MUX	00: Source IR[11:9] 01: Source IR[8:6] 10: Source R6
LD.MAR	0: load not enabled 1: load enabled	ALUMUX	0: Choose SR1 1: Choose TEMP
LD.MDR	0: load not enabled		
	1: load enabled	ALUK	00: ADD 01: AND
LD.CC	0: load not enabled		10: NOT
	1: load enabled		11: Pass input A
LD.TEMP	0: load not enabled	MIO.EN	0: MIO not enabled
	1: load enabled		1: MIO enabled
GatePC	0: do not pass signal	R.W	0: Read
	1: pass signal		1: Write
GateMDR	0: do not pass signal		
	1: pass signal		
GateALU	0: do not pass signal		
	1: pass signal		

Part b: What does the new instruction do?

Name:
Problem 9. (20 points): Consider a two player game where the players must think quickly each time it is their turn to make a move. Each player has a total allotted amount of time to make all his/her moves. Two clocks display the remaining time for each player. While a player is thinking of his/her move, his clock counts down. If time runs out, the other player wins. As soon as a player makes his/her move, he hits a button, which serves to stop counting down his clock and start counting down the other player's clock.
The program on the next page implements this mechanism. The main program keeps track of the time remaining for each player by decrementing the proper counter once per second while the player is thinking. When a player's counter reaches zero, a message is printed on the screen declaring the winner. When a player hits the button, an interrupt is taken. The interrupt service routine takes such action as to enable the main program (after returning from the interrupt) to start decrementing the other counter.
The interrupt vector for the button is $x35$. The priority level of the button is $#2$. Assume that the operating system has set the Interrupt Enable bit of the button to enable it to interrupt. Assume the main program runs at priority $#1$ and executes in user mode.
Part a : In order for the interrupt service routine to be executed when the button is pushed, what memory location must contain what value?
Address: Value:
Part b : Assume a player hits the button while the instruction at line 16 is being executed. What two values (in hex) will be pushed on the stack?
Part c: Fill in the missing instructions in the user program.
Part d : This program has a bug that will only occur if an interrupt is taken at an inappropriate time. Write down the line number of an instruction such that if the button is pressed while that instruction is executing, unintended behavior will result.
Line Number:
How could we fix this bug?

Name:_ ; Interrupt Service Routine .ORIG x1550 NOT R0, R0 RTI .END ; User Program .ORIG x3000 ; Line 1 AND R0, R0, #0 LDR1, TIME ; Line 2 LD R2, TIME ; Line 3 NEXT BRn P2_DEC ; Line 6 ADD R1, R1, #-1 ; Line 7 RO, P2WINS LEA ; Line 9 BRnzp END ; Line 10 P2_DEC ADD R2, R2, #-1 ; Line 11 RO, P1WINS LEA ; Line 13 END PUTS ; Line 14 ${\tt HALT}$; Line 15 COUNT LD R3, SECOND ; Line 16 LOOP ADD R3, R3, #-1 ; Line 17 BRp LOOP ; Line 18 #300 TIME .FILL

TIME .FILL #300
SECOND .FILL #25000 ; 1 second
P1WINS .STRINGZ "Player 1 Wins."
P2WINS .STRINGZ "Player 2 Wins."
.END

