Computer Architecture: Fundamentals, Tradeoffs, Challenges

Chapter 9: Input/Output (I/O)

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Outline

- Characteristics of I/O
- Bus Transactions
- An example: asynchronous bus with central arbitration
 - Arbitration
 - Race Conditions
 - Transfer
- RAID (Redundant Array of Disks)

Characteristics of I/O

Three parts

- The medium (e.g., the magnetic field in the track)
- The device itself (e.g., the disk)
- The controller
- How
 - Polling
 - Interrupt driven
 - DMA (the I/O control block)
 - I/O processor
- Instructions
 - Memory-mapped
 - Special I/O instructions

Bus Transactions

- Synch vs Asynch
 - Asynch (slow)
 - Handshaking
 - No clock
 - Everything explicit
 - Synch (fast)
 - Clock → Most things are implicit
 - Very fast, but must be short
- Signals
 - Three types: Address, Data, Control
 - Multiplexed address, data
- Arbitration
 - Central: Priority Arbitration Unit
 - Distributed: my "dinner table" analogy
- Transfer

An Asynchronous I/O System



Arbitration

• The concept:



• If the device does not want the bus:



Is there a problem?

A Race Condition

- Consider the following:
 - The PAU asserts the BG signal
 - Device A does not want the bus
 - Controller A passes it on
 - Controller B wants it, asserts SACK
 - Controller A sees SACK, returns to Idle
- What if:
 - Device A wants the bus before PAU negated BG
 - Controller A goes to BR and, since BG is still asserted
 - Controller A goes to SACK
- How do we fix it?

The fix!

- We do not return to IDLE when we see SACK
 PAU may still be asserting BG
- We wait until PAU stops asserting BG – Then it is safe to return to IDLE
- The fix:



What if a higher priority request comes in AFTER the PAU has issued BG?

How do we keep PAU from issuing higher BG
Disable new requests to PAU at start of bus cycle

(Bus master, asserts BBSY, negates SACK)

 Enable requests to PAU at end of arbitration (Next bus master asserts SACK)



The Transfer



Redundant Array of Interdependent Disks (RAID)

- The soul of RAID: performance plus redundancy
 - Introduced by Norman Ouchi (IBM), US Patent granted in 1978
 - Acronym by Gibson, Katz, Patterson (UC Berkeley), 1988
- The meaning of I in RAID
 - Initially Inexpensive, until they realized it was not inexpensive
 - Then independent, except the disks are not independent
 - I suggest "Interdependent" !
- The various levels
 - RAID 0: Vanilla -- Coarse, No redundancy
 - RAID 1: Mirroring Coarse, Redundancy
 - RAID 2: ECC Fine, ECC
 - RAID 3: Parity Fine, Parity disk
 - RAID 4: Coarse parity Coarse, Parity disk
 - RAID 5: The preferred model Fine, no parity disk
 - RAID 6: More than one mechanism for error checking

The RAID levels

- RAID 0:
 - Coarse
 - No Redundancy
- RAID 1:
 - Coarse
 - Redundancy





- RAID 2:
 - Fine

- ECC



The RAID levels (continued)

- RAID 3:
 - Fine
 - Parity Disk
- RAID 4:
 - Coarse
 - Parity Disk



- Fine
- No Parity Disk







Danke!