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## **Final Exam**

Date: December 15, 2017

Printed Name:	Last,	First	
_	your promise that you have not cheated and wi will not reveal the contents of this exam to othe		
Signature:			

## **Instructions:**

- Write your UT EID on all pages (at the top) and circle your instructor's name at the bottom.
- Closed book and closed notes. No books, no papers, no data sheets (other than the last four pages of this Exam)
- No devices other than pencil, pen, eraser (no calculators, no electronic devices), please turn cell phones off.
- Please be sure that your answers to all questions (and all supporting work that is required) are contained in the space (boxes) provided. *Anything outside the boxes will be ignored in grading*.
- You have 180 minutes, so allocate your time accordingly.
- For all questions, unless otherwise stated, find the most efficient (time, resources) solution.
- Unless otherwise stated, make all I/O accesses friendly.
- Please read the entire exam before starting. See supplement pages for Device I/O registers.

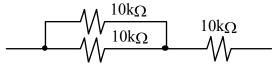
Problem 1	10	
Problem 2	10	
Problem 3	10	
Problem 4	15	
Problem 5	15	
Problem 6	15	
Problem 7	15	
Problem 8	10	
Total	100	

(10) Question 1. Place one letter in the box for each question that represents the best answer.  A) To force the compiler to not optimize the access  B) To place the variable in nonvolatile ROM  C) To force the variable to be placed in a register  D) To place the variable in volatile RAM  E) To make the variable private to the file  F) To make the variable private to the function  G) Because of Arm Architecture Procedure Call Standard  H) To force the variable to be placed on the stack	
(1) Part a) Why do we add static to an otherwise local variable?	
(1) Part b) Why do we add const to an otherwise global variable?	
(1) Part c) Why do we add static to an otherwise global variable?	
(1) Part d) Why do we use a local variable?	
I) To execute instructions faster J) Because of the Central Limit Theorem K) To reduce the latency of other interrupts L) To save power, making the battery last longer M) To prevent Aliasing N) Because of the Nyquist Theorem O) To decouple the execution of the ISR with the main program P) To reduce noise and improve signal to noise ratio Q) To synchronize one computer to another R) To improve bandwidth	
(1) Part e) Why do we place a FIFO queue between an ISR that reads data from an input port	
(1) Part f) Why does a Harvard architecture have two (or more) buses?	
(1) Part g) Why should the time to execute an ISR be as short as possible?	
(1) Part h) Why would you ever wish to use the PLL and slow down the bus clock	
ı	
(1) Part i) Why would we use hardware averaging on the ADC?	
(1) Part j) Why does the UART protocol use start and stop bits?	

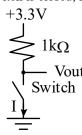
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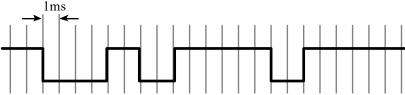
- (10) Question 2: For parts a) to f), please place one numerical value in the box for each question.
- (2) Part a) You are measuring heart rate using a bioelectric signal and need frequency components from 1 to 10 Hz. The ADC precision is 8 bits. UART baud rate is 100,000 bps. What is the ......slowest sampling rate (sampling frequency) possible?
- (1) Part c) The bit time is 1ms. What is the baud rate?
- (1) Part e) These three resistors could be replaced by one resistor. Give the resistance value of this one equivalent resistor?



(1) Part f) How much current (I) flows when the switch is closed, include units ......



Reverse-engineer UART parameters from the trace observed at a receiver below.



- (2) Part g) What is the *data value* transferred over the UART in **unsigned decimal**? .....
- (2) Part h) What is the baud rate in bits/sec? .....

that each	a) Compi bus cycl	lete the assembly subroutine is 62.5 ns. The interrupt of 3600 seconds. You may as	the that initializes SysTick to interrupt every 1 sec. The bus clock is 16 MHz; occurs when the counter goes from 1 to 0. Initialize CTRL to 7. The goal is to sume PB1 is already initialized to be a GPIO output. Fill in the blanks as
	. text		
Init:	LDR	R1,=SysTick_LOAD	
	STR	R0,[R1]	// establish interrupt period
	LDR	R1,=SysTick_CTRL	
	STR	R2,[R1]	// arm SysTick interrupts
			//clear bit 0 of PRIMASK

BX LR

**(4) Part b)** Write the SysTick ISR in assembly that toggles PB1 every 3600 seconds. Write 2 to GPIOB\_DOUTTGL31\_0 to toggle PB1. You cannot read GPIOB\_DOUTTGL31\_0.

SysTick_Handler:		

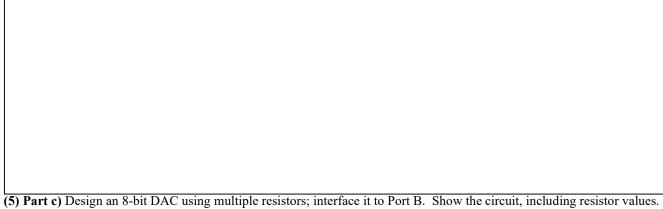
(5) Part a) Draw a Moore finite state graph with 2 inputs (SW1, SW0) and one output (LOCK). Initially, the LOCK is hig (1). Any time both inputs are simultaneously high, the machine reverts to the initial state with the LOCK high. Call this initial state Init. The sequence just SW0 high (input=1), both switches low (input=0), and then just SW1 high (input=2) will cause the LOCK to go low (0). All other states have the LOCK high. Assume the FSM runs at a fixed rate of 100 Hz periodically performing output, wait 10ms, input, and next operations. Use pointer addressing to access the next state.	
(10) Part b) Show the C code, including the struct that defines your finite state graph in ROM. Define a state pointer Pt.	
(10) 1 are by one or the code, including the struct that defines your finite state graph in Novi. Define a state pointer 1 t.	
(10) 1 are b) onon the code, including the state that defines your finite state graph in ROM. Define a state pointer 11.	
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(10) Face by Show the C code, including the state that defines your finite state graph in NOW. Define a state pointer 11.	
(10) Fare by Show the C code, including the strate that defines your finite state graph in NOW. Define a state pointer II.	
(10) Fart b) Show the C code, including the state that defines your filme state graph in KOWI. Beline a state pointer I t.	
(10) 1 are by onon, the e code, including the strate that defines your finite state graph in NOW. Define a state pointer I t.	

The execution engine in the main program is given and cannot be changed.

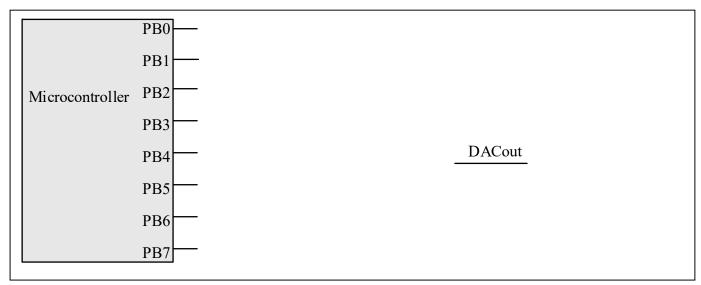
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۱	13	ĮV	uesuon	J.	mer	jucing

(5) Part a) Interface an LED to the microcontroller PA7 output. If PA7 is high the LED should be on. Assume the desired operating point is 2.1V 2mA. Assume the output low voltage of the MSPM0 is 0.2V. Assume the output high voltage of the MSPM0 is 3.1V. Show the circuit, including resistor values.

(5) Part b) Design a circuit with two switches and one LED. Assume the switches are ideal, and the LED operating point is 1.3V 2 mA. There is no microcontroller in this solution. The LED should be on if both switches are pressed; otherwise, the LED should be off. Show the circuit, including the switches, the LED, and resistors as needed (include resistance values).



Give actual resistor values that will work.



- (15) Question 7: Local variables
- (3) Part a) What does the function Func do?

(7) Part b) Implement the operation performed by this assembly function in C. The assembly is written in AAPCS. Use good names for parameters and local variable(s).

```
//R0, pt points to an array
//R1, the size of array
Func: PUSH {R0,R1}
     MOVS R2,#0
                   ;result
     PUSH {R2,R7}
    MOV R7,SP
//***this point for part c)***
loop: LDR R0,[R7,#8]
     LDRH R1, [R0]
     ADDS R0,#2
     STR R0, [R7,#8]
     LDRH R0, [R7,#0]
     EORS R0,R0,R1
     STRH R0, [R7,#0]
     LDRB R0, [R7, #12]
     SUBS R0,#1
     STRB R0, [R7, #12]
     BNE loop
     LDRH R0, [R7,#0]
     ADD SP,#4
     POP
         {R7}
     ADD
          SP,#8
     BX
          LR
```

(5) Part c) Assume at the beginning of the execution of the function Func, the stack is empty, SP equals 0x20200400, R0 equals 0x20201000 (pointer to an array), R1 equals 4 (size of the array), R2 equals 2 (some random irrelevant value), and R7 equals 11 (some random irrelevant value). Show the contents of the stack at the point in the execution of Func signified by the \*\*\*. Also specify the value of R7 by drawing an arrow into the stack figure.

0x202003E4	
0x202003E8	
0x202003EC	
0x202003F0	
0x202003F4	
0x202003F8	
0x202003FC	
0x20200400	
0x20200404	

(10) Question 8: <u>UART interrupt</u>

Assume you have the SSI interface to the LCD you used in Lab 5. In C, this LCD output data function is

```
void SPI_OutData(char data) {
  while((SPI1->STAT&0x02) == 0x00) {}; // spin if TxFifo full
  GPIOA->DOUTSET31_0 = 1<<13; // RS=PA13=1 for data
  SPI1->TXDATA = data;
}
```

The overall goal of the communication is to transfer all data from the UART0 receiver to the LCD. Assume the UART0 is initialized for receiver interrupts. When there is data in the receiver FIFO, RXRIS is set and a UART interrupt is triggered.

(6) Part a) Show the UART ISR that reads one 8-bit data from the receiver, acknowledges the interrupt and sends the data to the LCD using the SSI port. To read data, you read from UART0->RXDATA.

- (2) Part b) The hardware automatically pushes R0, R1, R2, R3, R12, LR, PC, and PSW on the stack when the interrupt is triggered. These registers are automatically popped at the end of the ISR by the BX LR instruction. Is the ISR code allowed to use the other registers R4-R11? Choose the best answer, placing A-F in the box.
- A) No, since these registers are not saved, using these registers would cause errors in the main program.
- B) Yes, according to AAPCS, R4-R11 are reserved for interrupts, so they can be freely used.
- C) Yes, according to AAPCS, any function can use R4-R11 if it first saves the registers and then restores them afterward.
- D) No, according to AAPCS, R4-R11 are reserved for the main program, so they cannot be used in an ISR.
- E) No, according to AAPCS, the stack must be aligned to 8 bytes.
- F) Yes, according to AAPCS, R4-R11 are automatically saved by the compiler for every function.
- (1) Part c) Does the hardware automatically disable interrupts during the execution of the ISR? Choose the best answer, placing A-F in the box.
- A) Yes, the execution of an interrupt is more important than the execution of the main program.
- B) Yes, disabling interrupts prevents the execution of one ISR from being interrupted by itself
- C) Yes, according to AAPCS, the hardware automatically disables interrupts during the execution of the ISRs.
- D) No, but according to AAPCS, the software automatically disables interrupts during the execution of the ISRs.
- E) No, interrupts are not disabled to allow interrupts with a lower priority (higher value in priority register) to interrupt.
- F) No, interrupts are not disabled to allow interrupts with a higher priority (lower value in priority register) to interrupt.